SCALING LIMITS OF RANDOM TREES AND GRAPHS

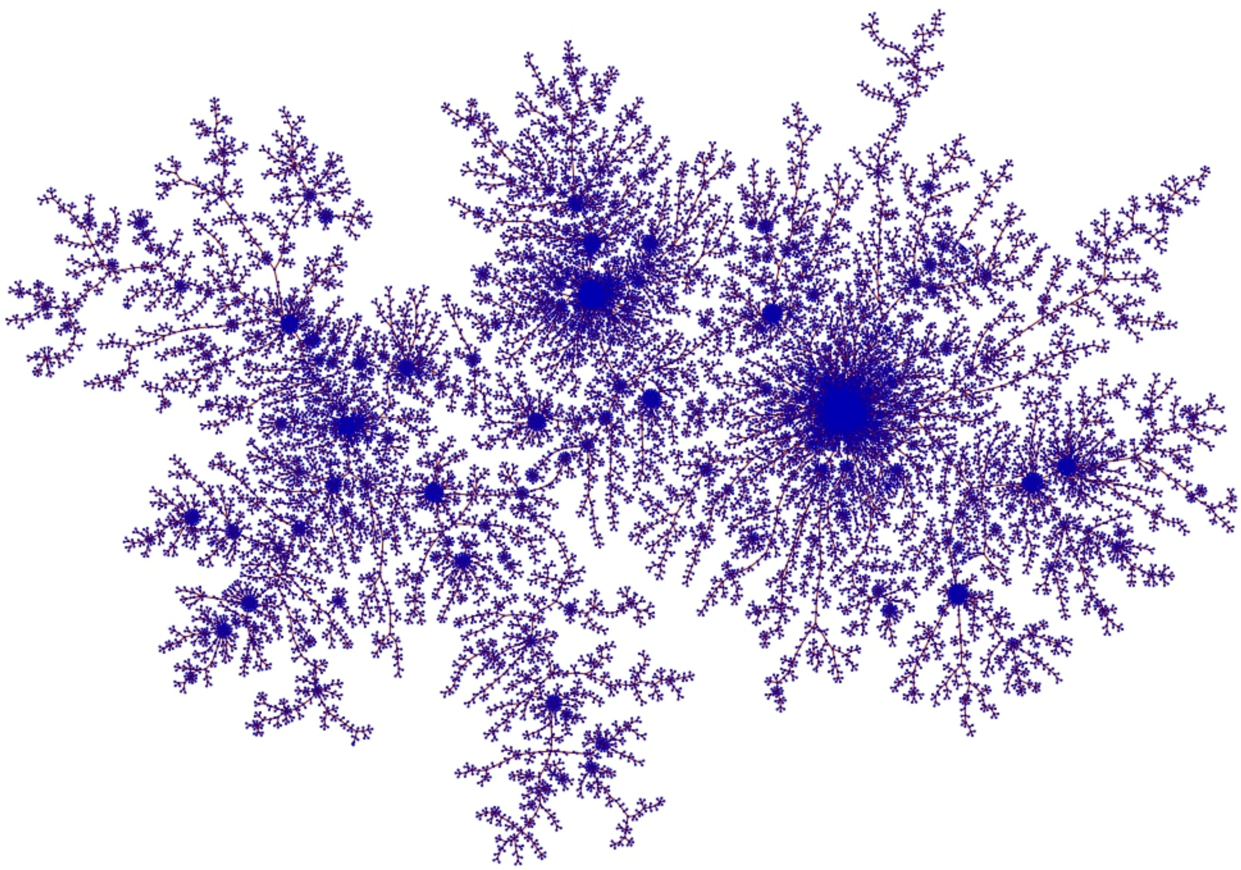

Figure 1: An image of a cool tree stolen from Igor Kortchemski's website.

Contents

Ramblings for people who stumble upon this file

The other day I woke up to a notifcation that google scholar added these notes to my profile. I guess this means people might actually end up reading these notes, so I think I should add some remarks about what these notes even are.

These notes were made for an informal course on scaling limits of random graphs at McGill in the winter 2025 semester (as it is winter 2025 right now, expect a lot of typos and half-baked ideas that were covered better in lecture than here). The intention of the course was to bring graduate students researching combinatorial probability theory up to speed with both the classical and modern work on scaling limits for random trees and graphs. Focus was placed on introducing and proving the results from metric geometry and probability theory that pre-date the ideas of graph scaling limits and supported the emergence of it. Much of the content from the first three sections was developed by expanding upon the excellent introduction to scaling limits provided in [LG05].

Thank you to the many attendees of the course who gave me a reason to actually learn this material well enough to present it. Special thanks in particular go to my PhD supervisors Luc Devroye and Louigi Addario-Berry for helping me out with the preparation and presentation of the material.

1 RANDOM COMBINATORIAL TREES

This section introduces our main object of consideration, which is random trees. We discuss two ways to encode trees with discrete functions and examine the relationships between these encodings. We then turn our attention to random trees, where the specific trees of interest are Bienaymé trees.

1.1 ENCODING TREES WITH DISCRETE FUNCTIONS

Most trees we consider in these notes are *plane trees*, which are finite rooted trees with an ordering on each collection of siblings in the tree. We shall identify all plane trees as subsets of the infinite Ulam-Harris tree, which we define now. Let

$$\mathbf{U} = \bigcup_{k=0}^{\infty} \mathbb{N}^k,$$

where we take $\mathbb{N} = \{1, 2, 3, ..., \}$ and $\mathbb{N}^0 = \{\emptyset\}$. We call the elements of $\mathbf{U}$ the *vertices*. The length of the vector $u \in \mathbf{U}$, $|u|$, is called the *generation* of u. It is also called the *height* of u. If $u = (u_1, ..., u_k), v = (v_1, ..., v_m) \in \mathbf{U}$ we let $uv = u \cdot v$ denote the concatenation of the two sequences, $(u_1, ..., u_k, v_1, ..., v_m)$. The vertex $p(v) = (u_1, ..., u_{k-1})$ is called the *parent* of u and u is called the child of $p(u)$. If $w = (w_1, ..., w_k) \in \mathbf{U}$ is such that $w_i = u_i$ for all $1 \leq i \leq k-1$ and $w_k \neq u_k$, then u and w are called *siblings*. The set $\mathbf{U}$ is called the *Ulam-Harris tree* (Figure 2 highlights the tree structure), and we use it to formally define the notion of a plane tree.

Definition 1.1. *A finite subset* $\mathbf{t} \subseteq \mathbf{U}$ *is called a* plane tree *if:*

(i) $\emptyset \in \mathbf{t}$.

(ii) *If* $u \in \mathbf{t}$*, then* $p(u) \in \mathbf{t}$.

(iii) *There is a collection of non-negative integers* $(c_{\mathbf{t}}(u) : u \in \mathbf{t})$ *such that, for all* $j \in \mathbb{N}$ *and* $u \in \mathbf{t}$*,* $uj \in \mathbf{t}$ *if and only if* $1 \leq j \leq c_{\mathbf{t}}(u)$.

We interpret $c_{\mathbf{t}}(u)$ as the number of children that u has in $\mathbf{t}$. We also occasionally refer to this as the *out-degree* of u. The set of all plane trees is denoted by $\mathcal{R}$ in what follows. The set of all plane trees $\mathbf{t}$ such that $|\mathbf{t}| = n$ is denoted by $\mathcal{R}_n$. The ordering on our plane trees is the natural lexicographical ordering of the Ulam-Harris tree. We shall occasionally need to discuss the genealogical partial ordering of our trees as well, which we shall denote with $\preceq$. We write $u \preceq v$ for two vertices $u, v \in \mathbf{t}$ if v is a descendent of u, i.e., $v = uw$ for some $w \in \mathbf{U}$. The lexicographical ordering of $\mathbf{U}$ is denoted with $\leq$.

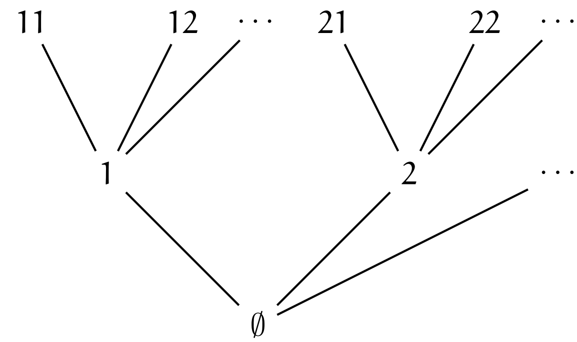

Figure 2: A depiction of the set **U** that highlights its tree structure.

The embedding of our plane trees inside the Ulam-Harris tree, and the corresponding ordering, allow for easy exploration of the tree via depth-first exploration. We first define the depth-first queue process, and then note why it is useful for characterizing plane trees.

Definition 1.2. *Let* $\mathbf{t} \in \mathcal{R}_n$ *and let* $u_1, ..., u_n$ *be the vertices written in lexicographical order. Write* $(c_1, ..., c_n) = (c_{\mathbf{t}}(u_1), ..., c_{\mathbf{t}}(u_n))$*. The sequence of integers* $(q_k)_{k=0}^n$ *with*

$$q_k = \sum_{i=1}^{k} (c_i - 1)$$

is called the depth-first queue process *of the tree* $\mathbf{t}$ *(DFQ). Any sequence* $(x_k)_{k=0}^n$ *such that*

(i) $x_0 = 0$*,* $x_n = -1$*,*

(ii) $x_k \geq 0$ *for all* $0 \leq k \leq n-1$

(iii) $x_k - x_{k-1} \geq -1$ *for all* $1 \leq k \leq n$

is called a Łukasiewicz path of length n*. We take* $\mathcal{L}$ *to denote the collection of all Łukasiewicz paths and* $\mathcal{L}_n$ *the paths of length* n*. In some places, the DFQ process of a tree is called the Łukasiewicz path of the tree.*

As the name suggests, there is an interpretation of the DFQ process of a tree $\mathbf{t} \in \mathcal{R}_n$ as the evolving size of a queue while exploring the tree. Begin with a queue $Q_0 = (\emptyset)$. Then, for $0 \leq i \leq n-1$, suppose that $Q_i = (w_1, ..., w_{q_{i+1}})$ with $q_i = |Q_i| - 1$. We pop w_1 from Q_i, query the number of children it has, and then add those children to the front of Q_i in their lexicographical order to form Q_{i+1}. The net change in the size of the queue at each step is exactly $c_i - 1$, as at each step the vertex being popped is the ith in the ordering of $\mathbf{t}$. Note that step k of the DFQ process is when we explore the vertex u_k (the kth vertex in the lexicographical order) and its children are not represented in the queue until the next step if it has any. Starting the walk at zero and not one is just a notational choice to make future convergence results a little cleaner. It removes a lot of "+1's".'

Lemma 1.3. *The mapping* $\varphi : \mathcal{R} \to \mathcal{L}$ *given by*

$$\varphi(\mathbf{t}) = (q_0, ..., q_{|\mathbf{t}|}) \quad \forall \mathbf{t} \in \mathcal{R},$$

where $(q_0, ..., q_{|\mathbf{t}|})$ *is the DFQ process for* $\mathbf{t}$*, is a bijection.*

Proof. First, we verify that φ maps into $\mathcal{L}$, which amounts to showing (i) and (ii) in the definition as the other point is clear. The first point follows from the fact that trees on n vertices have $n-1$ edges (and hence $n-1$ children in the context of plane trees). For the second point, we note that $c_{\mathbf{t}}(u_1) + ... + c_{\mathbf{t}}(u_k) \geq k$ for $1 \leq k \leq n-1$ because $u_1, ..., u_{k+1}$ are all children of some vertex in $\{u_1, ..., u_k\}$.

Recall that two plane trees $\mathbf{t}, \mathbf{s}$ are equal if and only if they are the same subset of $\mathbf{U}$. We begin by showing that φ is injective. If $|\mathbf{t}| \neq |\mathbf{s}|$, then they do not have the same DFQ process so suppose that $|\mathbf{t}| = |\mathbf{s}| = n$ and $\mathbf{t} \neq \mathbf{s}$. Let $u^* \in \mathbf{t} \cap \mathbf{s}$ be the first vertex in the ordering that has a child in one tree and not the other. Without loss of generality, we may assume that this child is in $\mathbf{t}$, so $c_{\mathbf{t}}(u^*) > c_{\mathbf{s}}(u^*)$. If $(q_0(\mathbf{t}), ..., q_n(\mathbf{t}))$ and $(q_0(\mathbf{s}), ..., q_n(\mathbf{s}))$ are the DFQ processes of $\mathbf{t}$ and $\mathbf{s}$ respectively, the fact that u^* was chosen to be minimal implies that $q_k(\mathbf{t}) = q_k(\mathbf{s})$ for all $1 \leq k \leq i^* - 1$, where i^* is the place of u^* in the ordering. Then,

$$q_{i^*}(\mathbf{t}) = q_{i^*-1}(\mathbf{t}) + c_{\mathbf{t}}(u^*) > q_{i^*-1}(\mathbf{s}) + c_{\mathbf{s}}(u^*) = q_{i^*}(\mathbf{s}).$$

Surjectivity follows almost immediately from the fact that $q_k - q_{k-1} = c_{\mathbf{t}}(u_k) - 1$ for all $1 \leq k \leq n$. Given a Łukasiewicz path $\mathbf{q} = (q_0, ..., q_n)$ we can construct a tree that straightforwardly maps to $\mathbf{q}$. Begin with $\mathbf{t}_0 = \{\emptyset\}$. Then, inductively define $\mathbf{t}_{i+1}$ for each $0 \leq i \leq n-1$ by setting $\mathbf{t}_{i+1} = \mathbf{t}_i \cup \{x_i \cdot 1, ..., x_i \cdot (q_{i+1} - q_i + 1)\}$, where x_i is the ith element of $\mathbf{t}_i$ in lexicographical order (note that such an element exists by the assumption $q_k \geq 0$ for $0 \leq k \leq n-1$). One can check that $\varphi(\mathbf{t}_n) = (q_0, ..., q_n)$. $\square$

Another discrete function that encodes plane trees is the height function. It can be seen as a walk through the tree in lexicographical order that records the height of the current vertex.

Definition 1.4. *Let* $\mathbf{t} \in \mathcal{R}_n$ *and let* $u_0, ...u_{n-1}$ *be its vertices written in lexicographical order. The* height function *of* $\mathbf{t}$*, denoted by* $(h_{\mathbf{t}}(k))_{k=0}^{n-1}$*, is given by* $h_{\mathbf{t}}(k) = |u_{k+1}|$.

Before we get into why the height function acutally matters, let's first introduce a continuous function that is related to the height function and of great importance later on. We call this function the *contour function* of the tree. The formal definition is a little confusing, I recommend looking at the example below to make sense out of it. We informally can see the contour function as arising from a process where we trace out the tree using a pencil that never leaves the paper and draws at a single unit speed. When we deal with the contour function we often take an intuitive approach, arguing with pictures and words instead of dealing with the formal objects. This just helps us to avoid long detours with a lot of notation that end with us concluding relatively intuitive statements. Anyways, here is the definition.

Definition 1.5. *Let* $\mathbf{t} \in \mathcal{R}_n$ *and let* $u_0, ..., u_{n-1}$ *be the vertices in lexicographical order. Set* $u_n = \emptyset$*. Let* $p_0^i, p_1^i, p_2^i...$ *be the interior vertices on the unique paths from* u_i *to* u_{i+1} *for each* $0 \le i \le n-1$ *in the order they would be taken if travelling from* u_i *to* u_{i+1} *in* $\mathbf{t}$*. We define a new sequence of vertices* $v_0, ..., v_{2(n-1)}$ *by inserting the* p^i*'s between* u_i *and* u_{i+1} *for all* $0 \le i \le n-1$ *(each vertex* $u \in \mathbf{t}$ *appears* $c_{\mathbf{t}}(u) + 1$ *times in the new sequence). We define the contour function of* $\mathbf{t}$*,* $\gamma_{\mathbf{t}} : [0, \infty) \to [0, \infty)$ *by*

$$\gamma(t) = |v_{\lfloor t \rfloor}| + (t - \lfloor t \rfloor)(|v_{\lceil t \rceil}| - |v_{\lfloor t \rfloor}|)$$

for $0 \le t \le 2(n-1)$*, and* $\gamma(t) = 0$ *for* $t > 2(n-1)$*.*

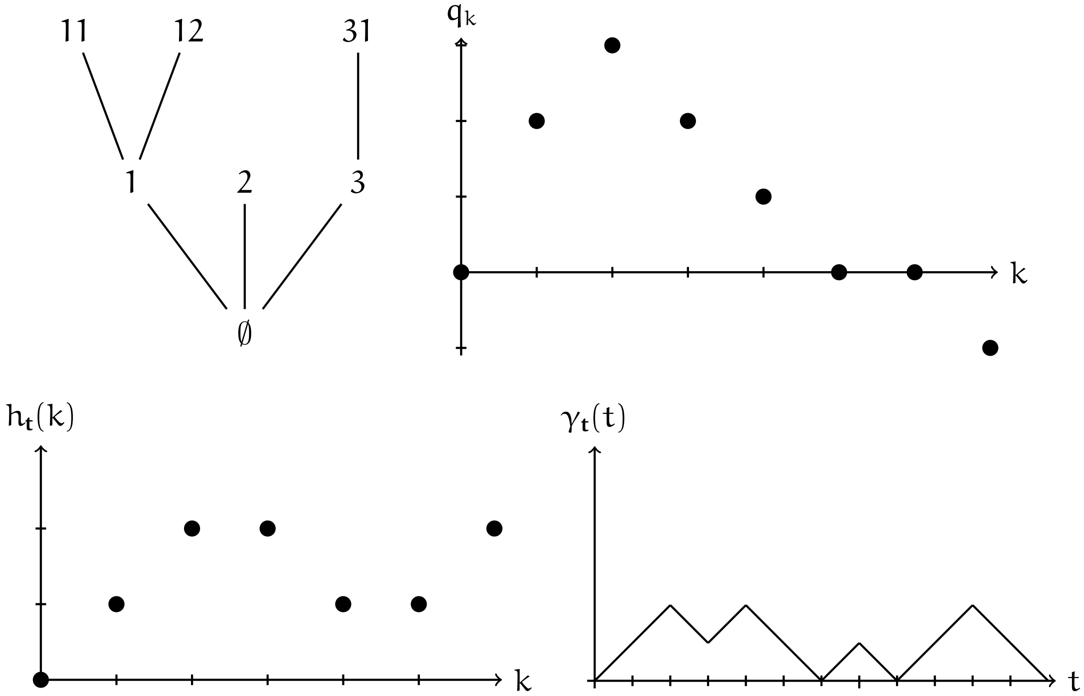

Figure 3: A tree and its many functional encodings

There is a simple way to convert between the height function and the DFQ process of a tree. This relationship will allow us to describe the height function in terms of sums of i.i.d. random variables when discussing Bienaymé trees later.

Theorem 1.6. *Let* $\mathbf{t} \in \mathbb{R}_n$ *have DFQ process* $(q_0, ..., q_n)$*. Then, for all* $0 \le k \le n-1$*,*

$$h_{\mathbf{t}}(k) = \left| \left\{ 1 \le j \le k-1 : q_j = \inf_{j \le m \le k} q_m \right\} \right|.$$

Proof sketch. It is clear that $h_t(k) = |\{0 \le j \le k-1 : u_j \preceq u_k\}|$, so we only need to show that

$$u_j \preceq u_k \iff q_j = \inf_{j \le m \le k} q_m.$$

It can be observed immediately from the definition that, if $\mathbf{t}(u_j)$ is the subtree of $\mathbf{t}$ rooted at u_j, then $u_j \preceq u_k$ if and only if $u_k \in \mathbf{t}(u_j)$, so we can instead show

$$u_k \in \mathbf{t}(u_j) \iff q_j = \inf_{j \le m \le k} q_m. \tag{1}$$

Let $\tau_j = \inf\{m \ge j : q_m < q_j\}$. At step j of the DFQ process we add u_j's children to the queue and remove u_j. The process only leaves the subtree $\mathbf{t}(u_j)$ all of the children of u_j have been removed (along with any children they have). This is exactly τ_j. In particular, we have that $\mathbf{t}(u_j) = \{u_m : j \le m \le \tau_j - 1\}$. (1) follows immediately from this identity. □

A corollary of Theorem 1.6 is that the height function of a tree uniquely determines it. By taking the end point of all length one intervals on which the contour function is increasing, we can recover the height process of a tree. Moreover, from the height function we can recover the tree and from the tree we can get the contour function. Hence, the contour function uniquely determines the tree as well. Of course, one can prove this fact directly via the "pencil and paper" analogy. One can also prove the height function encodes its tree directly by observing that, if one knows the u_k and $h_{\mathbf{t}}(k+1)$, then there is only one possible vertex that could be u_{k+1} (it is a child of the ancestor of u_k that is at height $h_{\mathbf{t}}(k+1)-1$). I'm being a bit handwavy here, but the conclusion really is just that all three of the processes presented here uniquely determine our trees.

1.2 Bienaymé trees

Definition 1.7. *Let μ be a measure on $\mathbb{Z}_{\ge} = \{0, 1, 2, ...\}$ with $\sum_{k=0}^{\infty} k\mu(k) < \infty$ such that $\mu(1) \neq 1$. For all $u \in \mathbf{U}$, we associate an independent random variable $\xi_u \overset{\mathcal{L}}{=} \mu$. The subset $T = \{u = (u^1, ..., u^k) \in \mathbf{U} : u^j \le \xi_{(u^1,...,u^{j-1})}\ \forall\ 1 \le j \le k\}$ is called a* Bienaymé tree *with offspring distribution μ. We often write $T \overset{\mathcal{L}}{=}$ Bienaymé(μ). Collections of many i.i.d. Bienaymé trees are sometimes called Bienaymé forests. We call a Bienaymé tree critical if $\sum_{k=0}^{\infty} k\mu(k) = 1$, subcritical if $\sum_{k=0}^{\infty} k\mu(k) < 1$, and supercritical otherwise.*

These trees are ubiquitous in probability theory and combinatorics, having been studied as far back as the 1800's. Those familiar with the classic Galton-Watson martingale process may notice that these two structures are essentially the same. It is mostly straightforward to prove from the definition that Bienymé trees are plane trees except for the criteria that T must be finite. This fact is a corollary of a result known by many as the fundamental theorem of Bienaymé trees. See [ANN04] for a proof.

Theorem 1.8. *Let $T \overset{\mathcal{L}}{=}$ Bienaymé(μ) for some μ matching the above criteria. If T is sub-critical or critical, then $|T| < \infty$ almost surely. In particular, T is a plane tree. Otherwise, $\mathbf{P}(|T| = \infty) > 0$.*

The independence in the variables $(\xi_u : u \in \mathbf{U})$ has some nice consequences concerning the distribution of T over the set $\mathcal{R}$.

Lemma 1.9. *Let $\mathbf{t} \in \mathcal{R}$ and let $T \stackrel{\mathcal{L}}{=}$ Bienaymé(μ). Then,*

$$\mathbf{P}(T = \mathbf{t}) = \prod_{u \in \mathbf{t}} \mu(c_{\mathbf{t}}(u)).$$

Proof. Since T is a plane tree almost surely, $\{T = \mathbf{t}\} = \cap_{u \in \mathbf{t}} \{\xi_u = c_{\mathbf{t}}(u)\}$. Using the independence of the ξ's we get,

$$\mathbf{P}(T = \mathbf{t}) = \mathbf{P}\left(\bigcap_{u \in \mathbf{t}} \{\xi_u = c_{\mathbf{t}}(u)\}\right) = \prod_{u \in \mathbf{t}} \mu(c_{\mathbf{t}}(u)).$$

□

With the standard pleasantries out of the way, we can turn our attention to the most important property of Bienaymé trees from the perspective of scaling limits. The DFQ process of these trees is distributed like a simple random walk, and their sizes are exactly distributed like the first time that the simple random walk hits -1. At first glance, knowing the definition of the DFQ process, one might think that this statement is trivially true by the definition of Bienaymé trees. However, the presence of the stopping time in the expression below makes the claim not immediate as it could (in theory) disturb the natural independence between the number of children each vertex has.

Theorem 1.10. *Let $T \stackrel{\mathcal{L}}{=}$ Bienaymé(μ), and let its DFQ process be denoted by Q. Let $(S_k : k \geq 0)$ be a simple random walk with step sizes distributed like ν, where for all $k \geq -1$, $\nu(k) = \mu(k+1)$. Then,*

$$Q \stackrel{\mathcal{L}}{=} (S_0, \ldots, S_\tau),$$

where $\tau = \inf\{n \geq 1 : S_n = -1\}$. In particular $|T| \stackrel{\mathcal{L}}{=} \tau$.

Proof. It suffices to just check that the vector $(c_{\mathbf{t}}(U_0), \ldots, c_{\mathbf{t}}(U_{|T|-1}))$ is distributed like a collection of i.i.d. μ-distributed random variables, where $(U_0, \ldots, U_{|T|-1})$ is the vertices of T written in lexicographical order. To be able to remove the random indexing, we want $\{U_k = u\}$ for $0 \leq k \leq |T| - 1$ and $u \in \mathcal{U}$ to be measurable with respect to only the vertices below u in the lexicographical order.

First, the set $T \cap \{v \in \mathbf{U} : v \leq u\}$, is measurable with respect to $\sigma(\xi_v : v < u)$. Then, for any $k \geq 0$, the event $\{U_k = u\} \cap \{|T| > k\}$, being completely determined by $T \cap \{v \in \mathbf{U} : v < u\}$, is measurable with respect to $\sigma(\xi_v : v < u)$. The set $\{U_k = u\} \cap \{|T| \leq k\}$ is also measurable with respect to $\sigma(\xi_v : v < u)$ for the same reason. Combining the two facts we get that $\{U_k = u\}$ is measurable with respect to $\sigma(\xi_v : v < u)$.

Now, from here we can proceed via a standard induction. Let $g_0, ..., g_k : \mathbb{Z}_\geq \to \mathbb{Z}_\geq$ be a collection of functions for $0 \leq k \leq |T| - 1$. Then,

$$\begin{aligned}
&\mathbf{E}\left[g_1(\xi_{U_0})\cdots g_k(\xi_{U_k})\right] \\
&= \sum_{u_0<\ldots<u_k} \mathbf{E}\left[\mathbf{1}_{\{U_0=u_0,\ldots,U_k=u_k\}} g_1(\xi_{u_1})\cdots g_k(\xi_{u_k})\right] \\
&= \sum_{u_0<\ldots<u_k} \mathbf{E}\left[\mathbf{1}_{\{U_0=u_0,\ldots,U_k=u_k\}} g_1(\xi_{u_1})\cdots g_{k-1}(\xi_{u_{k-1}})\right] \mathbf{E}[g_k(\xi_{u_k})] \\
&= \sum_{u_0<\ldots<u_{k-1}} \mathbf{E}\left[\mathbf{1}_{\{U_0=u_0,\ldots,U_{k-1}=u_{k-1}\}} g_1(\xi_{u_1})\cdots g_{k-1}(\xi_{u_{k-1}})\right] \mathbf{E}[g_k(\xi_{u_k})] \\
&= \mathbf{E}\left[g_1(\xi_{U_0})\cdots g_k(\xi_{U_{k-1}})\right] \mathbf{E}[g_k(\xi_{u_0})],
\end{aligned}$$

where in the first equality we used the measurability we just proved and in the second we use the independence of child distribution for fixed indices. The sum is only over vertices in generation at most k. Applying induction completes the proof of the independence, and as noted at the start completes the proof as a whole. □

1.3 Bienaymé tree conditioned to have a fixed size

Bienaymé trees are interesting structures in their standard form. However, their ability to generalize so many canonical random tree models is what has kept them an ongoing topic of discussion for so many years since their origins in the study of family trees. The way we observe this generalizing property is by sampling Bienaymé trees conditioned on their size being some parameter $n \in \mathbb{N}$. We write $T \overset{\mathcal{L}}{=} \text{Bienaymé}(n, \mu)$ for a random plane tree T if, for all $\mathbf{t} \in \mathcal{R}_n$,

$$\mathbf{P}(T = \mathbf{t}) = \mathbf{P}(T' = \mathbf{t} \mid |T'| = n),$$

where $T' \overset{\mathcal{L}}{=} \text{Bienaymé}(\mu)$. For the rest of this subsection, we are going to cover a variety of random tree models, and explain how they fit into the category of conditioned critical Bienaymé trees. First, however, we need to explain why this is something that we should be able to do.

Definition 1.11. *Let M be a multiset of plane trees. We define the weight of a tree in $\mathbf{t} \in \mathbf{U}$, $\Omega(\mathbf{t})$, to be the number of occurrences of $\mathbf{t}$ in M. Then, we call*

$$z_n = \sum_{\mathbf{t} \in M : |\mathbf{t}| = n} \Omega(\mathbf{t})$$

the partition function of M. For each $n \geq 1$, let T_n be a random tree with distribution,

$$\mathbf{P}(T_n = \mathbf{t}) = \frac{\Omega(\mathbf{t})}{z_n}.$$

For each $\mathbf{t} \in \mathcal{U}$, let $(m_k(\mathbf{t}))_{k=0}^{\infty}$ be the number of vertices with k children for $k \geq 0$. If there exists a sequence $(a_k)_{k=1}^{\infty}$ of integers such that

$$\Omega(\mathbf{t}) = \prod_{k=0}^{\infty} a_k^{m_k(\mathbf{t})},$$

then we call the random trees $(T_n)_{n=1}^{\infty}$ a simply generated family of random trees.

In many cases, simply generated trees can be described as Bienaymé trees conditioned on their size. Let $(T_n)_{n=1}^{\infty}$ be a family of simply generated tree, and let μ^x be a measure defined by $\mu^x(k) = a_k x^k / f(x)$ for all $k \geq 0$ and some $x > 0$. We define T_n^x for all $n \geq 1$ to be a Bienaymé(n, μ^x).

Lemma 1.12. *Let $f(x) = \sum_{k=0}^{\infty} a_k x^k$ and suppose that there is some $x^* > 0$ such that $1 \leq f(x^*) < \infty$. Then, there exists some $\tau > 0$ such that $f(\tau) = \tau f'(\tau)$.*

We shall skip the proof as it not particularly instructive and generating functions are not the topic of interest.

Theorem 1.13. *Let $f(x) = \sum_{k=0}^{\infty} a_k x^k$ and suppose that there is some $x^* > 0$ such that $1 \leq f(x^*) < \infty$. Let $\tau > 0$ such that $f(\tau) = \tau f'(\tau)$ (exists from the above lemma). Then, for all $x \in (0, \tau]$, $T_n \stackrel{\mathcal{L}}{=} T_n^x$, where both $(T_n)_{n=1}^{\infty}$ and $(T_n^x)_{n=1}^{\infty}$ are defined above. In particular, there is a critical child distribution μ such that $T_n \stackrel{\mathcal{L}}{=}$ Bienaymé(n, μ).*

Proof. Let $T^* \stackrel{\mathcal{L}}{=}$ Bienaymé(μ^t). By Lemma 1.9,

$$\begin{aligned}
\mathbf{P}(T^* = \mathbf{t}) &= \prod_{k=0}^{\infty} (\mu^x(k))^{m_k(\mathbf{t})} \\
&= \prod_{k=0}^{\infty} \left(\frac{a_k x^k}{f(x)} \right)^{m_k(\mathbf{t})} \\
&= \left(\prod_{k=0}^{\infty} a_k^{m_k(\mathbf{t})} \right) (f(x))^{-n} \left(x^{\sum_{k=0}^{\infty} k m_k(\mathbf{t})} \right) \\
&= \Omega(\mathbf{t})(f(x))^{-n} \left(x^{\sum_{k=0}^{\infty} k m_k(\mathbf{t})} \right).
\end{aligned}$$

Then,

$$\mathbf{P}(|T^*| = n) = \sum_{\mathbf{t}:|\mathbf{t}|=n} \Omega(\mathbf{t})(f(x))^{-n} \left(x^{\sum_{k=0}^{\infty} k m_k(\mathbf{t})} \right) = z_n (f(x))^{-n} \left(x^{\sum_{k=0}^{\infty} k m_k(\mathbf{t})} \right).$$

Hence,

$$\mathbf{P}(T_n^x = \mathbf{t}) = \frac{\Omega(\mathbf{t})}{z_n}.$$

The second statement follows the above lemma and the fact that the mean of the child distribution μ^x is

$$\sum_{k=0}^{\infty} \frac{k a_k x^k}{f(x)} = \frac{x f'(x)}{f(x)}.$$

□

What is the takeaway of this theorem? Our claim at the beginning of this section was that we could view many canonical random tree models as Bienaymé trees conditioned on their size. This theorem just asserts that we only need to be able to view them as simply generated trees, which is a much nicer family for this purpose. It is fairly easy to find a weight function that results in the correct distribution for many families of random trees. Let us finish things off by giving some examples. Verifying the claims is not too hard and I don't even know if I'll cover this material, so I'm just going to write the coefficients that give the desired tree for each example.

(i) If we set $a_0 = 1,\ a_1 = 2,\ a_2 = 1$, then T_n is a uniform rooted binary tree on n vertices.

(ii) If we set $(a_0 = 1,\ a_2 = 1)$, then T_n is a uniform full binary tree on n vertices.

(iii) If we set $(a_0 = 1,\ a_k = 1)$, then T_n is a uniform rooted k-ary tree on n vertices.

(iv) If we set $(a_k = 1$ for all $k \geq 0)$, then T_n is a uniform rooted plane tree on n vertices.

There is one last case that needs to be separated out on its own as we can deal directly with the Bieanymé tree instead of the simply generated tree. The tree of interest is the uniform random labelled tree on n vertices. Let $T \overset{\mathcal{L}}{=} \text{Bienaymé}(\text{Poi}(1))$. Erase the planar ordering and root, and then give T a uniformly chosen labelling from $\{1, ..., |T|\}$. Then, for a labelled rooted tree $\mathbf{t}$,

$$\mathbf{P}(T = \mathbf{t}) = \frac{e^{-|\mathbf{t}|}}{|\mathbf{t}|!},$$

implying that $\mathbf{P}(T = \mathbf{t} \mid |T| = n)$ is a uniform labelled tree on n vertices (the identity is not trivial, but can be verified without too much sweat by permuting vertices with the same degree).

2 Real trees and the Brownian CRT

We introduce a second notion of a tree in this section, specifically that of a real tree. These are connected metric spaces that share metric information with combinatorial trees, but erase the meaning of things like vertices and adjacency. We discuss how the space of all real trees can be made into a complete separable metric space, setting ourselves up the groundwork for how one can make sense out of scaling limits for trees. We also cover the encoding of real trees via continuous functions supported on a compact connected set. This sets up a bridge between the combinatorial and the continuum via the contour function.

2.1 The space of rooted real trees

As was done with combinatorial trees, we shall begin our exploration of real trees by setting them up as formal structures. Naturally, the starting place is the definition.

Definition 2.1. *A compact metric space* $(\mathbf{T}, d)$ *is called a real tree if, for all* $x, y \in \mathbf{T}$*:*

(i) *there is a unique is isometric embedding* $f_{xy} : [0, d(x, y)] \to \mathbf{T}$ *such that* $f_{xy}(0) = x$ *and* $f_{xy}(d(x, y)) = y$*;*

(ii) *if* $g : [0, 1] \to \mathbf{T}$ *is a continuous injective map with* $g(0) = x$ *and* $g(1) = y$*, then* $g([0, 1]) = f([0, d(x, y)])$.

Despite no longer feeling like vertices in the sense that they are in a combinatorial tree, we shall still call elements of $\mathbf{T}$ its *vertices*. The real trees we discuss in these notes shall be rooted, meaning that each $\mathbf{T}$ has some distinguished vertex $\rho \in \mathbf{T}$. Its role shall mostly be as a constraint for the equivalence of two trees, though its existence also allows to discuss things like height. Real trees are not considered planar, but some results we prove later about how much branching can occur in a real tree imply that we could define an ordering analogous to the sibling ordering that defines plane trees. We need some more notation to go along with our new definition.

(i) The range of the isometric embedding f_{xy} for any $x, y \in \mathbf{T}$ shall be denoted by $[x, y]$. The sets $(x, y]$, $[x, y)$, (x, y), $[x, x]$, $(x, x]$, $[x, x)$, (x, x) are all defined analogously.

(ii) The distance $d(\rho, x)$ for $x \in \mathbf{T}$ is called the *height* of x. The segment $[\rho, x]$ is called the ancestral line of x.

(iii) We define the *genealogical partial ordering* on $\mathbf{T}$, written as $\preceq$, by $x \preceq y$ if $x \in [\rho, y]$.

(iv) The *degree* of a vertex $x \in \mathbf{T}$ is the cardinality of the set of components in the metric space $(\mathbf{T} \setminus \{x\}, d)$. We say that y and z are in the same component of $\mathbf{T} \setminus \{x\}$ if they are connected in $\mathbf{T} \setminus \{x\}$ in the topological sense. Vertices of degree one are called *leaves*.

(v) For $x, y \in \mathbf{T}$, we call the unique $z \in \mathbf{T}$ such that $[\rho, x] \cap [\rho, y] = [\rho, z]$ the *least common ancestor* of x and y. We denote this vertex by $x \wedge y$.

(vi) We call two real trees $\mathbf{T}_1$ and $\mathbf{T}_2$ *equivalent* if there is a root preserving isometry $f : \mathbf{T}_1 \to \mathbf{T}_2$. The set $\mathbb{T}$ will denote the space of all equivalence classes of real trees. We often conflate a tree with its equivalence class.

Item (v) above contained the claim that there exists such an element. Since it gives us a chance to get acquainted with the definition of a real tree, let's prove this claim.

Lemma 2.2. *For every pair* $x, y \in \mathbf{T}$*, there exists a unique vertex* $z \in \mathbf{T}$ *such that* $[\rho, x] \cap [\rho, y] = [\rho, z]$.

Proof. Let $a = \sup\{b \in [0, d(\rho, x)] : f_{\rho x}(b) \in [\rho, y]\}$, and let $z = f_{\rho x}(a)$. By the closeness of the sets $[\rho, x]$ and $[\rho, y]$, we know that $z \in [\rho, x] \cap [\rho, y]$, implying that $[\rho, z] \subseteq [\rho, x] \cap [\rho, y]$. On the other hand, if $z' \in [\rho, x] \cap [\rho, y]$, then $f_{\rho x}^{-1}(z') \in \{b \in [0, d(\rho, x)] : f_{\rho x}(b) \in [\rho, y]\}$, and so $f_{\rho x}^{-1}(z') \leq a$. Using the fact that $f_{\rho x}$ is an isometric embedding we can see that $d(\rho, z) = a$ and that $f|_{[0,a]}$ is the unique isometric embedding of $[0, d(\rho, z)]$ into $\mathbf{T}$. Hence, $z' \in [\rho, z]$ and $[\rho, x] \cap [\rho, y] \subseteq [\rho, z]$. Uniqueness is straightforward. If $[\rho, x] = [\rho, y]$ for any $x, y \in \mathbf{T}$, then $x \preceq y$ and $y \preceq x$. In particular $x = y$. □

There are many equivalent notions of real trees. Almost all of them use (i) (which is called the unique geodesic condition), but (ii) (the no-loop property) could be restated in any number of ways [Jan23]. Item (i) also is the property that asserts connectedness. There is one common equivalent description that does not use (i) and we shall record it because it is fun. Rather than pretend that I can say anything about the proof, I shall simply state it and bask in its glory ([Jan23] discusses this equivalent definition as well if you would like to learn about it).

Theorem 2.3. *A compact rooted metric space* (X, d) *is a real tree if and only if it is path-connected and satisfies the four-point condition :*

$$d(x_1, x_2) + d(x_3, x_4) \leq \max\{d(x_1, x_3) + d(x_2, x_4), d(x_1, x_4) + d(x_2, x_3)\},$$

for all $x_1, x_2, x_3, x_4 \in X$.

Ok, moving on. With the goal of convergence theorems in mind, we would like to have a notion of distance between two real trees. In most cases, our particular choice of distance function is the Gromov-Hausdorff distance. There are multiple equivalent definitions of this distance,and we take the following one to be our canonical definition. For $(\mathbf{T}_1, d_1)$ and $(\mathbf{T}_2, d_2)$ real trees, we call $C \subseteq \mathbf{T}_1 \times \mathbf{T}_2$ a (root-preserving) correspondence between $\mathbf{T}_1$ and $\mathbf{T}_2$ if:

(i) $\forall x_1 \in \mathbf{T}_1 \; \exists x_2 \in \mathbf{T}_2$ such that $(x_1, x_2) \in C$,

(ii) $\forall x_2 \in \mathbf{T}_2 \; \exists x_1 \in \mathbf{T}_1$ such that $(x_1, x_2) \in C$, and

(iii) $(\rho_1, \rho_2) \in C$, where ρ_1 and ρ_2 are the roots of the trees $\mathbf{T}_1$ and $\mathbf{T}_2$ respectively.

The space of all correspondences between $\mathbf{T}_1$ and $\mathbf{T}_2$ is denoted by $\mathcal{C}(\mathbf{T}_1, \mathbf{T}_2)$. Then, we define the Gromov-Hausdorff distance between $(\mathbf{T}_1, d_1)$ and $(\mathbf{T}_2, d_2)$ as

$$d_{GH}(\mathbf{T}_1, \mathbf{T}_2) = \frac{1}{2} \inf_{C \in \mathcal{C}(\mathbf{T}_1, \mathbf{T}_2)} \operatorname{dis}(C),$$

where

$$\operatorname{dis}(C) = \sup \left\{ |d_1(x_1, y_1) - d_2(x_2, y_2)| : (x_1, x_2),\ (y_1, y_2) \in C \right\}.$$

There is a slightly more intuitive definition of the GH distance in terms of the Hausdorff distance of isometric embeddings of $\mathbf{T}_1$ and $\mathbf{T}_2$ into a mutual space. This definition will be of use later down the line, and for this sake we introduce it now.

Definition 2.4. *The Hausdorff distance* d_H *between two compact sets* K_1, K_2 *of a metric space* (X, d) *is defined by*

$$\inf\{\epsilon > 0 : K_1 \subseteq K_2^\epsilon, K_2 \subseteq K_1^\epsilon\},$$

where $S^\epsilon = \{x \in X : d(x, S) \leq \epsilon\}$.

Lemma 2.5. *For two real trees* $(\mathbf{T}_1, d_1)$ *and* $\mathbf{T}_2, d_2)$ *with roots* ρ_1 *and* ρ_2 *we define a metric*

$$d(\mathbf{T}_1, \mathbf{T}_2) = \inf_{\varphi_1, \varphi_2} \left(d_H(\varphi(\mathbf{T}_1), \varphi(\mathbf{T}_2)) \vee d^*(\varphi_1(\rho_1), \varphi_2(\rho_2)) \right),$$

where the infimum is taken over all isometric embeddings of $\mathbf{T}_1$ *and* $\mathbf{T}_2$ *and choices of destination* (X^*, d^*).

Proof. First, suppose that $d(\mathbf{T}_1, \mathbf{T}_2) < r$ for two trees $(\mathbf{T}_1, d_1)$ and $(\mathbf{T}_2, d_2)$ and let φ_1, φ_2 be isometric embeddings into a space (Z, d_Z) such that $d_H(\varphi_1 \mathbf{T}_1, \varphi_2 \mathbf{T}_2) < r$. We define a relation C by adding all pairs of vertices $(t_1, t_2) \in \mathbf{T}_1 \times \mathbf{T}_2$ such that $d_Z(\varphi_1(t_1), \varphi_2(t_2)) < r$. By the assumption at the beginning, C is a correspondance that with $\operatorname{dis}(C) < 2r$. To see this, consider two pairs of corresponding points (x_1, x_2)

and (y_1, y_2), and suppose that $d_1(x_1, y_1) \geq d_2(x_2, y_2)$. Then, a simple application of the triangle inequality gives

$$\begin{aligned}
&d_1(x_1, y_1) - d_2(x_2, y_2) \\
=&d_Z(\varphi_1 x_1, \varphi_1 y_1) - d_Z(\varphi_2 x_2, \varphi_2 y_2) \\
\leq&d_Z(\varphi_1 x_1, \varphi_2 x_2) + d_Z(\varphi_2 x_2, \varphi_1 y_1) - d_Z(\varphi_2 x_2, \varphi_2 y_2) \\
\leq&d_Z(\varphi_1 x_1, \varphi_2 x_2) + d_Z(\varphi_2 x_2, \varphi_1 y_2) + d_Z(\varphi_2 y_2, \varphi_1 y_1) - d_Z(\varphi_2 x_2, \varphi_2 y_2) \\
=&d_Z(\varphi_1 x_1, \varphi_2 x_2) + d_Z(\varphi_2 y_2, \varphi_1 y_1),
\end{aligned}$$

which is strictly below $2r$ by definition. Hence, we can conclude that $d_{GH} \leq d$. Now suppose that $\mathrm{dis}(C) = 2r$ for some correspondance C. Then, in the disjoint union of $\mathbf{T}_1$ and $\mathbf{T}_2$ (mark all the points in $\mathbf{T}_1$ with a zero and in $\mathbf{T}_2$ with a one and then take the union) we define a pseudometric

$$d^*(t_1, t_2) = \begin{cases} \inf_{(t_1', t_2') \in C} \Big(d_1(t_1, t_1') + d_2(t_2, t_2') + r \Big), \text{ if } t_1 \in \mathbf{T}_1, t_2 \in \mathbf{T}_2 \\ d_1(t_1, t_2), \text{ if } t_1, t_2 \in \mathbf{T}_1 \\ d_2(t_1, t_2), \text{ if } t_1, t_2 \in \mathbf{T}_2 \end{cases}.$$

Note that $d^*(t_1, t_2) = r$ when the two vertices correspond with each other. In particular, since every vertex has a partner in the correspondance (and the roots correspond), we have that $d_H(\mathbf{T}_1, \mathbf{T}_2) \leq r$. There are some issues with the fact that d^* is only a pseudometric, but simply modding out by the standad distance zero equivalence relation finishes the job. □

An important remark to make is that there was nothing special about the fact that our compact metric spaces of choice were trees in any of the proof of any of those definitions. One can extend the notion of Gromov-Hausdorff distance that we just provided to the set of all isometry classes of compact metric spaces. We will often make reference to this larger space containing $\mathbb{T}$ when working with real trees and especially when working with real graphs. We denote it by $\mathbb{K}$. The last thing to cover about Gromov-Hausdorff space before moving on to functional encodings is the question of completeness.

Theorem 2.6. *Both* $(\mathbb{K}, d_{GH})$ *and* $(\mathbb{T}, d_{GH})$ *are complete separable metric spaces.*

Proof sketch. Separability of $\mathbb{K}$ is not too hard to show with the correspondance definition of the Gromov-Hausdorff distance. Since our metric spaces are compact, we can find finite ϵ-covers of them for all $\epsilon > 0$. This implies that the set of finite metric spaces is dense in $\mathbb{K}$. If we take all finite metric spaces that have only rational distances, then we get a countable dense set. We can do a very similar thing for $\mathbb{T}$ by considering all real trees that branch out only finitely many times

I didn't quite have time to type up a full argument for this proof following what was done in class. Hopefully I can fill this in later when I have time.

□

Due mostly to time constraints we have not ventured very deep into the theory of Gromov-Hausdorff space, only presenting the results that are needed. I would just like to remark that this is not due to lack of relevance or because the connetions end with what has been discussed here. Deep knowledge of the theory of convergence for metric spaces and the surrounding material has and will continue to be important to developing the theory of graph scaling limits. I recommend taking a look at [Bur01] to learn more about the topic, it was my main source of deeper information about Gromov-Hausdorff convergence when preparing these notes. I also stole a couple ideas from [Pet06].

2.2 Encoding real trees with functions

In this subsection, we argue why we can replace the study of real trees with the study of certain types of continuous functions. As noted in the summary of this section, this offers a bridge between the real trees of this section, and the plane trees of the previous section. First we set up our candidates for the encodings.

Let $f \in \{g : [0, \infty) \to [0, \infty) : \text{supp}(f) \text{ compact and connected}, g(0) = 0\} := C_c^+[0, \infty)$. We shall construct a real tree from the function. Define, for all $s, t \geq 0$,

$$m_f(s,t) = \inf_{\min(s,t) \leq r \leq \max(s,t)} f(r),$$

and $d_f(s,t) = f(s) + f(t) - 2m_f(s,t)$. Then, d_f is a metric on the set of equivalence classes $[0, \infty)/R_f$, where $R_f = \{(s,t) \in [0, \infty) \times [0, \infty) : d_f(s,t) = 0\}$. Essentially, our main theorem of this subsection asserts that the collection of all metric spaces $([0, \infty)/R_f, d_f)$ for functions $f \in C_c^+[0, \infty)$ is a rich enough set to fill our tree related needs. For a function $f \in C_c^+[0\infty)$, we let $(\mathbf{T}_f, df_f)$ denote the space $([0, \infty)/R_f, d_f)$ with root $\rho = [0]_{R_f}$, the equivalence class of 0 under R_f. It is relatively straightforward to show that $\mathbf{T}_f$ is in fact a compact metric space using uniform continuity of continuous functions over compact intervals, however we need to still show that they are real trees. In particulat, we would like the following to be true:

(i) For any $f \in C_c^+[0, \infty)$, the pair $(\mathbf{T}_f, d_f)$ is a real tree.

(ii) For any two real trees $(\mathbf{T}_f, d_f)$ and $(\mathbf{T}_g, d_g)$, $d_{GH}(\mathbf{T}_f, \mathbf{T}_g) = \Theta(\|f - g\|_\infty)$.

(iii) For every real tree $(\mathbf{T}, d)$, there exists a function $f \in C_c^+[0, \infty)$ such that $(\mathbf{T}, d) = (\mathbf{T}_f, d_f)$.

One way to show (i) is to observe that any metric spaces of the form $(\mathbf{T}_f, d_f)$ satisfy the four-point condition, which implies they are all real trees via Theorem 2.3. We shall take a more elementary approach that relies most on basic analysis techniques. To prove (i) and (ii) we first prove the results for almost-linear functions (defined below) and then invoke the completeness of $(\mathbb{T}, d_{GH})$ to extend to all functions in $C_c^+[0, \infty)$. A different approach to prove the same results that argues directly with the definition of a real tree is covered in [LG05]. The third point is actually not relevant in these notes and so we won't prove it. However, for the sake of completeing the analogy with the results from the previous section we think it is worthy to mention that (iii) is also true. An excellent constructive proof can be found in [Duq06].

Let $f \in C_c^+[0, \infty)$. We say that f is a almost-linear if there is $\epsilon, \Delta > 0$ such that for any $n \geq 0$ $f(x) = f(n\epsilon) + \Delta(x - n\epsilon)$ or $f(x) = f(n\epsilon) - \Delta(x - n\epsilon)$ for $x \in [n\epsilon, (n+1)\epsilon]$. We shall label the set of almost-linear functions in $C_c^+[0\infty)$ with C_L. We begin by asserting that almost linear-function produce real trees. We can conclude this fact by observing that the metric spaces produced by almost-linear functions are essentially equivalent to combinatorial plane trees.

Lemma 2.7. *Let $f \in C_L$. Then, the function $\gamma : [0, \infty) \to [0, \infty)$ given by $\gamma(t) = \Delta^{-1} f(\epsilon t)$ is the contour function for some plane tree $\mathbf{t}_f$. Moreover, $(\mathbf{T}_f, d_f)$ is isometric to the real tree version of $\mathbf{t}_f$ with edge lengths Δ.*

We shall skip past proving Lemma 2.7 or producing a formal construction of the real tree version of $\mathbf{t}_f$ with edge lengths Δ, favouring an appeal to intuition (see figure below). The idea is essentially that, as we sketch out the contour function with our pencil and paper, we can graft on intervals of length Δ every time that we begin an up interval for the function. One other thing worth observing is that ϵ actually plays no role in the structure of $(\mathbf{T}_f, d_f)$. This is not an issue and makes sense for what we want our functional encodings to be. We can see straight from the definition of $\mathbf{T}_f$ that, if we define $g(x) = f(\alpha x)$ for any $\alpha > 0$, the mapping $x \mapsto \alpha x$ induces an isometry $\mathbf{T}_f \to \mathbf{T}_g$.

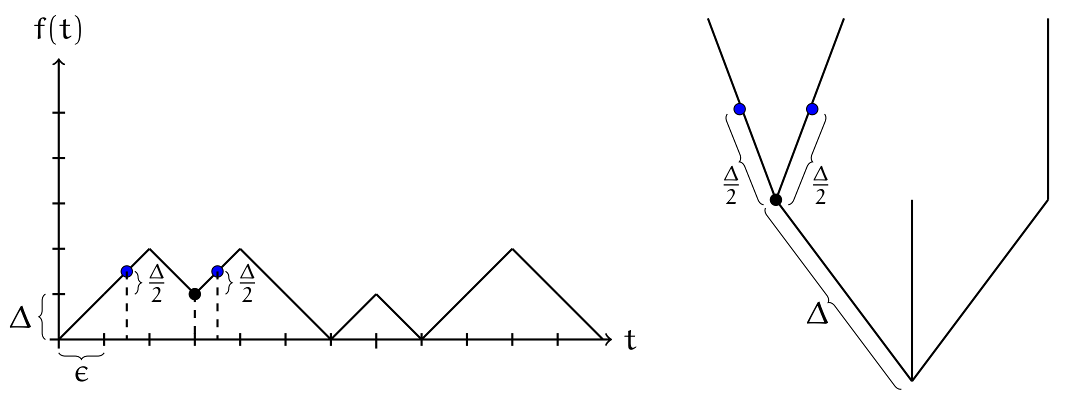

Figure 4: An almost-linear function and its corresponding real tree. Two points on the graph of the function are highlighted in blue, along with their corresponding vertices in the real tree to highlight how the distance d_f matches the natural extension of graph distance we get from sketching out the contour function. The greatest common ancestor of the points/vertices is black in both drawings.

Lemma 2.7 covers point (i) for the case of almost-linear functions. What is left to do is to argue that we can approximate all of the metric spaces for $C_c^+[0, \infty)$ via those generated by functions in C_L.

Lemma 2.8. C_L *is dense in* $C_c^+[0,\infty)$ *under the norm* $\|\cdot\|_\infty$.

Proof. It suffices to show the result for Lipschitz functions in $C_c^+[0,\infty)$ as they are dense in the set $C_c^+[0,\infty)$. Let $f \in C_c^+[0,\infty)$ be C-Lipschitz. Let $\Delta_n = C$ and $\epsilon_n = (S-I)n^{-1}$, where $S = \sup\operatorname{supp}(f)$ and $I = \inf\operatorname{supp}(f)$. Define recursively

$$P_n(j) = \begin{cases} +1, & \text{if } f(j\epsilon + I) \geq f_n(j\epsilon + I) \\ -1, & \text{otherwise} \end{cases}.$$

Finally, we set

$$f_n(t) = \sum_{j=0}^{(n-1)} P_n(j)\Delta_n\big((t-j\epsilon)_+ \vee \epsilon\big) - \sum_{j=0}^{f_n(S)(\Delta_n\epsilon_n)^{-1}} \Delta_n\big((t-S)-j\epsilon) \vee \epsilon\big).$$

The second sum exists only to make sure that the function is in $C_c^+[0,\infty)$ as promised, it disappears in the limit. We claim that $\|f - f_n\|_\infty \leq 2\Delta_n\epsilon_n$. We can proceed via induction. Suppose that $\sup_{x\in[I,k\epsilon+I]}|f_n(x) - f(x)| \leq 2\Delta_n\epsilon_n$ for some $0 \leq k < n-1$. Then, in particular $|f_n(k\epsilon+I) - f(k\epsilon+I)| \leq 2\Delta_n\epsilon_n$. There are two cases to consider. case 1: $f(k\epsilon + I) \geq f_n(k\epsilon + I)$. In this case the function f_n increases on the next interval. Since $|f(t) - f(k\epsilon + I)| \leq C(t - k\epsilon - I)$, we have that

$$\sup_{t\in[k\epsilon+I,(k+1)\epsilon+I]}(f(t)-f_n(t)) \leq f(k\epsilon+I)+C(t-k\epsilon-I)-f_n(k\epsilon+I)-C(t-k\epsilon-I) \leq 2\Delta_n\epsilon_n,$$

and

$$\sup_{t\in[k\epsilon+I,(k+1)\epsilon+I]}(f_n(t) - f(t)) \leq f(k\epsilon + I) + \Delta_n\epsilon_n - f_n(k\epsilon + I) - (-\Delta_n\epsilon_n) \leq 2\Delta_n\epsilon_n.$$

In particular, we have using the assumption that $\sup_{x\in[I,(k+1)\epsilon+I]}|f_n(x)-f(x)| \leq 2\Delta_n\epsilon_n$. case 2: $f(k\epsilon+I) < f_n(k\epsilon+I)$. This case goes almost identically to the first case so we shall omit this. We note that this induction actually extends to include times above S without changing anything as the second sum defining $f_n(t)$ is only empty when $f_n(S) > 0 = f(S)$. Thus, the proof is done as $\Delta_n\epsilon_n \to 0$ as $n \to \infty$. □

Combining the previous lemmas we can conclude what we wanted to show.

Theorem 2.9. *The two claims stated at the beginning of the section hold.*

(i) *For any two real trees* $(\mathbf{T}_f, d_f)$ *and* $(\mathbf{T}_g, d_g)$, $d_{GH}(\mathbf{T}_f, \mathbf{T}_g) \leq 2\|f - g\|_\infty$.

(ii) *For any* $f \in C_c^+[0,\infty)$, *the pair* $(\mathbf{T}_f, d_f)$ *is a real tree.*

Proof. (i) can be proven using the correspondance definition of the Gromov-Hausdorff distance (we have not yet shown that the metric spaces are trees, but recall that we can define the GH-distance for any two compact metric spaces). Let

$$C = \big\{([x]_{R_f}, [y]_{R_g}) : \exists t \geq 0 \text{ such that } t \in [x]_{R_f},\ t \in [y]_{R_y}\big\}.$$

It can be observed easily that this is a root-preserving correspondence. Let $(x_1, y_1), (x_2, y_2) \in C$ (we are supressing the $[\cdot]_{R_f}$ now for clarity). Then, there exists $s, t \geq 0$ such that

$$|d_f(x_1, x_2) - d_g(y_1, y_2)| \leq |f(s) - g(s)| + |f(t) - g(t)| + 2|m_f(s,t) - m_g(s,t)|.$$

Without loss of generality we can assume that $m_f(s,t) \geq m_g(s,t)$. By the continuity of the two functions and the fact that $[s \wedge t, s \vee t]$ is closed there is some $p \geq 0$ such that $m_g(s,t) = g(p)$. Then,

$$2|m_f(s,t) - m_g(s,t)| \leq 2(f(p) - g(p)) \leq 2\|f - g\|_\infty.$$

Altogether, we get that

$$d_{GH}(\mathbf{T}_f, \mathbf{T}_g) \leq \frac{1}{2}\,\mathrm{dis}(C) \leq 2\|f - g\|_\infty.$$

We can easily prove (ii) using (i), the density of C_L in $C_c^+[0,\infty)$, and the fact that $\mathbb{T}$ is closed in $\mathbb{K}$. $\square$

3 Scaling limits of random walks and Bienaymé trees

We finally prove some scaling limits in this section. We begin with building up the theory of scaling limits for random functions, explaining the topological backing behind it and proving Donsker's Theorem. Using the theorem and results from the previous two sections, we prove scaling limits for the height function of both conditioned and un-conditioned critical Bienaymé trees. As a corollary, we obtain a scaling limit in the Gromov-Hausdorff topology for critical conditioned trees to a random real tree called the Brownian CRT. It is defined to be a real tree that is encoded by a unit length Brownian excursion.

3.1 Random functions in $C[0,1]$ and Donsker's Theorem

I borrowed a lot of the material in this subsection from [Bil13]. In order to discuss scaling limits, we require some results connecting random walks and Brownian motion. We also desire some good tools to explore the convergence of random functions with our functional encodings of real trees in mind. Our setup in this section is a sequence of i.i.d. random variables $(\xi_n)_{n\geq 1}$ with mean 0 and variance 1. Let $S_k = \sum_{i=1}^k \xi_i$. The sequence of random functions that we consider is $(W_n)_{n\geq 1}$, where $W_n : [0,1] \to \mathbb{R}$ is such that

$$W_n(t) = \frac{S_{\lfloor nt \rfloor} + (nt - \lfloor nt \rfloor)\xi_{\lceil nt \rceil}}{\sqrt{n}}. \tag{2}$$

Donsker's Theorem essentially asserts that the functions $W_n(t)$ converge towards Brownian motion on the interval $[0,1]$.

Theorem 3.1 (Donsker's Theorem)**.**

$$\Big(W_n(t) : t \in [0,1]\Big) \xrightarrow{\mathcal{L}} \Big(B(t) : t \in [0,1]\Big),$$

as $n \to \infty$ *in the space* $(C[0,1], \|\cdot\|_\infty)$*, where* $(B(t) : t \geq 0)$ *is standard one dimensional Brownian motion that starts with* $B(0) = 0$.

While we can intuitively view this theorem as being a sort of generalization of the central limit theorem (the sequence $(W_n(1)/\sqrt{n})_{n\geq 1}$ is exactly the sequence $(S_n/\sqrt{n})_{n\geq 1}$), we need to recall some topological tools to be able to complete the proof. This increased difficulty is due to the fact that the claimed convergence is in the space $C[0,1]$ rather than $\mathbb{R}$. Specifically, we desire an equivalence between convergence in distribution and convergence of finite dimensional marginals for continuous functions.

3.1.1 Convergence of measures on $C[0,1]$

Let us begin by dragging some old dusty theorems out from our attic.

Definition 3.2. *Let (X, τ) be a Hausdorff space and let $\mathcal{P}$ be the space of all probability measures on X equipped with the Borel sigma-algebra. A set $S \subseteq \mathcal{P}$ is called* tight *if for all $\epsilon > 0$ there is a compact set $K(\epsilon)$ such that $\sup_{\mu \in S} \mu(X \setminus K(\epsilon)) < \epsilon$.*

Theorem 3.3 (Prokhorov's Theorem)**.** *Let (X, d) be a separable metric space and let $\mathcal{P}$ be the set of all probability measures on X with the Borel sigma-algebra. Then, $S \subseteq \mathcal{P}$ is tight if and only if it is pre-compact.*

An almost direct consequence of Prokhorov's Theorem is worth recording.

Corollary 3.4. *Let $(\mu_n)_{n=1}^{\infty}$, μ be probability measures on $(C[0,1], \|\cdot\|_\infty)$. If the finite-dimensional marginals of $(\mu_n)_{n=1}^{\infty}$ converge in distribution to the finite-dimensional marginals of μ, and if $(\mu_n)_{n=1}^{\infty}$ is tight, then $\mu_n \xrightarrow{\mathcal{L}} \mu$ as $n \to \infty$.*

Proof. Recall that, for probability measures μ and ν on $[0,1]$, $\mu = \nu$ if and only if $\pi_{t_1,\ldots,t_k}\mu = \pi_{t_1,\ldots,t_k}\nu$ for $0 \leq t_1 \leq \ldots \leq t_k \leq 1$, where $\pi_{t_1,\ldots,t_k}$ is the projection onto the coordinates $t_1, \ldots, t_k$ (this can be observed by a standard $\pi - \lambda$ system proof).

Let $(\mu_{n_k})_{k=1}^{\infty})$ be a subsequence of $(\mu_n)_{n=1}^{\infty})$. By pre-compactness, this sequence has a convergent subsequence, tending to some limit μ^*. By the finite-dimensional marginals convergence and the fact from the previous paragraph, it holds that $\mu^* = \mu$. Hence, every subsequence of $(\mu_n)_{n=1}^{\infty}$ has a further subsequence that converges to μ. It is well known that this implies that $\mu_n \xrightarrow{\mathcal{L}} \mu$ as $n \to \infty$. □

Theorem 3.5 (Arzelà-Ascoli Theorem)**.** *A set $S \subseteq C[0,1]$ is pre-compact if and only if $\sup_{f \in S} |f(0)| < \infty$ and $\lim_{\delta \to 0} \sup_{f \in S} w_f(\delta) = 0$, where $w_f(\delta) = \sup_{|s-t|<\delta} |f(s) - f(t)|$ for all $0 < \delta < 1$.*

The function w_x is called the modulus of continuity. For our purposes, we need a translation of tightness into some criteria that are more easily verified by computations. We can begin by deriving a pair of conditions that mirror the pre-compactness definition given by the Arzelà-Ascoli Theorem.

Lemma 3.6. *A sequence of measures $(\mu_n)_{n=1}^{\infty}$ on $(C[0,1], \|\cdot\|_\infty)$ is tight if and only if the following two conditions hold:*

(i) *for all $\epsilon > 0$ there is $N, t \geq 0$ such that $\mu(\{x : |x(0)| > t\}) \leq \epsilon$ for all $n \geq N$,*

(ii) *for all $\epsilon > 0$, $\lim_{\delta \to 0} \limsup_{n \to \infty} \mu_n(\{x : w_x(\delta) \geq \epsilon\}) = 0$.*

Proof. Suppose that the sequence is tight. Choose some $K \subseteq C[0,1]$ and $t \geq 0$ such that $\mu_n(K) \geq 1 - \eta$. Then, by compactness, $K \subseteq \{x : |x(0)| \leq t\}$ and $K \subseteq \{x : w_x(\delta) \leq \epsilon\}$ for all $n \geq 1$ and $\delta > 0$ chosen sufficiently small by the Arzelà-Ascoli

Theorem. It quickly follows that $\mu_n(\{x : |x(0)| \geq t\}) \leq \eta$. Morover, we can see that $\lim_{\delta \to 0} \sup_{n \geq 1} \mu_n(\{x : w_x(\delta) \geq \epsilon\}) = 0$ by choosing K appropriately.

For the reverse direction, we may instead show the result under the assumption (ii)': for all $\eta, \epsilon > 0$ that $\mu_n(\{x : w_x(\delta) \geq \epsilon\}) \leq 1 - \eta$ for all n above some chosen $N \geq 0$.

Suppose that (i) and (ii)' hold for $N \geq 0$. We claim that each of the individual measures $\mu_1, ..., \mu_N$ are tight.

Since $C[0,1]$ is separable, we can find for each $k \geq 0$ a collection of balls of radius k^{-1}, $A_1, ..., A_{n_k}^{(k)}$ such that $\mu_1(\cup_{i=1}^{n_k} A_i^{(k)}) \geq 1 - \epsilon 2^{-k}$. The closure K of the set $\cap_{k=1}^{\infty} \cup_{i=1}^{n_k} A_i^{(k)}$ has measure $\mu_1(K) \geq 1 - \epsilon$ and is compact.

Returning back to the proof, a simple application of the union bound proves that the collection $\mu_1, ..., \mu_N$ is tight. This implies that the inequalities from (i) and (ii)' hold for this collection too. In particular, this allows us to assume that $N = 1$ in (i) and (ii)'. Choose some $t \geq 0$ such that $\mu_n(\{x : |x(0)| \leq t\}) \geq 1 - \epsilon$ for all $n \geq 1$ and choose δ_k such that $\mu_n(\{x : w_x(\delta_k) < k^{-1}\}) \geq 1 - \epsilon 2^{-k}$ for all $n \geq 1$. Then, if we set K to be the closure of

$$(\{x : |x(0)| \leq t\}) \cap \bigcap_{k=1}^{\infty} \left\{x : w_x(\delta_k) < k^{-1}\right\},$$

we have that $\mu_n(K) \geq 1 - 2\epsilon$ for all $n \geq 1$. By the Arzelà-Ascoli Theorem K is compact. □

In order to do probabilistic computations cleanly we need to be able to work with a nicer form of the modulus of continuity than is provided via its definition. Our final lemma covers this for us. Afterwards, we are left with criteria for weak convergence that are much more easily verified.

Lemma 3.7. *Suppose that* $0 = t_0 \leq ... \leq t_k = 1$ *is such that* $\min_{1 \leq i \leq k}(t_i - t_{i-1}) \geq \delta$. *Then, for any* $x \in C[0,1]$,

$$w_x(\delta) \leq 3 \max_{1 \leq i \leq k} \sup_{t_{i-1} \leq t \leq t_i} |x(t) - x(t_{i-1})|,$$

and

$$\mu(\{x : w_x(\delta) \geq 3\epsilon\}) \leq \sum_{i=1}^{k} \mu\left(\left\{x : \sup_{t_{i-1} \leq t \leq t_i} |x(t) - x(t_{i-1})| \geq \epsilon\right\}\right)$$

for any measure μ *on* $C[0,1]$.

Proof. The first inequality is a simple triangle inequality argument. Let

$$M = \max_{1 \leq i \leq k} \sup_{t_{i-1} \leq t \leq t_i} |x(t) - x(t_{i-1})|.$$

If $|s-t| \leq \delta$, then they are either in adjacent intervals or the same interval. Suppose that $s, t \in [t_{i-1}, t_i]$ for some chosen i. Then,

$$|x(s) - x(t)| \leq |x(s) - x(t_{i-1})| + |x(t) - x(t_{i-1})| \leq 2M.$$

Suppose that $s \in [t_{i-1}, t_i]$ and $t \in [t_i, t_{i+1}]$ for some chosen i. Then,

$$|x(s) - x(t)| \leq |x(s) - x(t_{i-1})| + |x(t_{i-1}) - x(t_i)| + |x(t) - x(t_i)| \leq 3M.$$

The second inequality follows from a union bound. □

3.1.2 Back to Donsker's Theorem

Equipped with Corollary 3.4, proving Donsker's Theorem is as easy as verifying the convergence for finite-dimensional marginals and the tightness condition.

Lemma 3.8. *Suppose that $(W_n)_{n=1}^{\infty}$ is defined as in (2). If*

$$\lim_{x\to\infty} \limsup_{n\to\infty} x^2 \mathbf{P}\left(\max_{1\leq k\leq n} |S_k| \geq x\sqrt{n}\right) = 0,$$

then the sequence $(W_n)_{n=1}^{\infty}$ is tight.

Proof. We proceed by showing the Arzelà-Ascoli conditions hold in Lemma 3.6. Condition (i) is immediate as $W_n(0) = 0$ for all $n \geq 1$, so we only need to verify the condition on the modulus of continuity for an arbitrary $\epsilon > 0$,

$$\lim_{\delta\to 0} \limsup_{n\to\infty} \mathbf{P}(w_x(W_n, \delta) \geq \epsilon) = 0.$$

Let $= m_0 \leq ... \leq m_k = n$, and consider times $t_i = \frac{m_i}{n}$. Applying Lemma 3.7 we get that

$$\mathbf{P}(w(W_n, \delta) \geq 3\epsilon) \leq \sum_{i=1}^{k} \mathbf{P}\left(\sup_{t_{i-1}\leq t\leq t_i} |W_n(t) - W_n(t_{i-1})| \geq \epsilon\right)$$

whenever $\delta \leq \frac{m_i - m_{i-1}}{n}$ for all $1 \leq i \leq k$. The chosen times are important because, by definition, $W_n(t_i) = S_{m_i}/\sqrt{n}$. Thus,

$$\sup_{t_{i-1}\leq t\leq t_i} |W_n(t) - W_n(t_{i-1})| = \frac{1}{\sqrt{n}} \max_{m_{i-1}\leq j\leq m_i} |S_j - S_{m_{i-1}}|,$$

and

$$\begin{aligned}
\mathbf{P}(w(W_n, \delta) \geq 3\epsilon) &\leq \sum_{i=1}^{k} \mathbf{P}\left(\frac{1}{\sqrt{n}} \max_{m_{i-1}\leq j\leq m_i} |S_j - S_{m_{i-1}}| \geq \epsilon\right) \\
&= \sum_{i=1}^{k} \mathbf{P}\left(\max_{0\leq j\leq m_i - m_{i-1}} |S_j| \geq \sqrt{n}\epsilon\right)
\end{aligned}$$

for appropriately chosen $(m_i)_{i=1}^k$ to suit the conditions on δ (the second equality is a consequence of the ξ_n's being i.i.d.). This bound leaves us with a much more familiar expression to deal with. First, we need to finalize our choices of parameters though.

Let $m = \lceil n\delta \rceil$, let $k = \lceil \delta^{-1} \rceil$, and let $m_i = 2im$ for each $0 \le i \le k$. Then, $m_i - m_{i-1} = m$ for all i and $(m_i - m_{i-1})/n \to 2\delta > \delta$ as $n \to \infty$.

With these chosen parameters the above expression becomes

$$\begin{aligned}\mathbf{P}(w(W_n, \delta) \ge 3\epsilon) &\le \delta^{-1}\mathbf{P}\left(\max_{0\le j\le 2m} |S_j| \ge \epsilon\sqrt{\frac{m}{\delta}}\right) \\ &= 2 \cdot (2\delta)^{-1}\mathbf{P}\left(\max_{0\le j\le 2m} |S_j| \ge \epsilon\frac{1}{\sqrt{2\delta}}\sqrt{2m}\right) \\ &= \frac{2}{\epsilon^2}x^2\mathbf{P}\left(\max_{0\le j\le 2m} |S_j| \ge x\sqrt{2m}\right),\end{aligned}$$

where we set $x = \epsilon(2\delta)^{-1/2}$. Note that, as $\delta \to 0$, $x \to \infty$. From here, applying the assumption is enough to yield condition (ii) in Lemma 3.6, which proves tightness. □

We are now ready to prove Donsker's Theorem, but first let us quickly recall the properties that characterize Brownian motion.

Definition 3.9. *One dimensional* Brownian motion *is a real-valued stochastic process* $(B(t) : t \ge 0)$ *that satisfies the following properties:*

(i) $B(0) = 0$.

(ii) *If* $t_0 < t_1 < \ldots < t_n$, *then* $B(t_0), B(t_1) - B(t_0), \ldots, B(t_n) - B(t_{n-1})$ *are independent.*

(iii) *If* $s < t$, *then* $B(s + t) - B(s) \stackrel{\mathcal{L}}{=} N(0, t - s)$.

These properties need to be shown for the limit of the finite-dimensional marginals of $(W_n)_{n=1}^\infty$ to complete the proof. It is enough to show that, for any collection of times $0 = t_0 \le \ldots \le t_k$ for some $k \ge 0$,

$$(W_n(t_1) - W_n(t_0), \ldots, W_n(t_k) - W_n(t_{k-1})) \xrightarrow{\mathcal{L}} (X_1, \ldots, X_k),$$

where the X_i's are independent with $X_i \stackrel{\mathcal{L}}{=} N(0, t_i - t_{i-1})$. This, along with tightness, is enough to complete the proof.

Theorem (Donsker's Theorem). *Let* $(\xi_n)_{n\ge1}$ *be a sequence of i.i.d. random variables with mean* 0 *and variance* 1. *Let* $S_k = \sum_{i=1}^k \xi_i$. *Define random functions* $(W_n)_{n\ge1}$ *where*

$$W_n(t) = \frac{S_{\lfloor nt \rfloor} + (nt - \lfloor nt \rfloor)\xi_{\lceil nt \rceil}}{\sqrt{n}}.$$

Then,

$$\Big(W_n(t) : t \in [0,1]\Big) \xrightarrow{\mathcal{L}} \Big(B(t) : t \in [0,1]\Big),$$

as $n \to \infty$ in the space $(C[0,1], \|\cdot\|_\infty)$, where $(B(t) : t \geq 0)$ is standard one dimensional Brownian motion that starts with $B(0) = 0$.

Proof. Let $t \geq s \geq 0$. $W_n(s) = S_{\lfloor ns \rfloor}/\sqrt{n} + X_n$ and $W_n(t) - W_n(s) = (S_{\lfloor nt \rfloor} - S_{\lfloor sn \rfloor})/\sqrt{n} + Y_n$, where X_n and Y_n are random variables that tend to 0 almost surely as $n \to \infty$. Basic properties of random walks assert that $S_{\lfloor ns \rfloor}$ and $(S_{\lfloor nt \rfloor} - S_{\lfloor sn \rfloor})$ are independent. By the central limits theorem and the continuous mapping theorem, we get that $W_n(s) \xrightarrow{\mathcal{L}} X$ and $W_n(t) - W_n(s) \xrightarrow{\mathcal{L}} Y$, where $X \stackrel{\mathcal{L}}{=} N(0,s)$ and $Y \stackrel{\mathcal{L}}{=} N(0, t-s)$ are independent. The general case is similar, and so we can move on to tightness. By Etemadi's inequality (see remark below if you are unfamiliar),

$$x^2 \mathbf{P}\left(\max_{0 \leq k \leq n} |S_k| \geq x\sqrt{n}\right) \leq 3x^2 \max_{0 \leq k \leq n} \mathbf{P}\left(|S_k| \geq x\sqrt{n}/3\right).$$

Let $k^*(x)$ be a constant depending on x chosen such that $\mathbf{P}(|S_k| \geq x\sqrt{k}/3) \leq \mathbf{P}(N(0,1) \geq x/3) + x^{-3}$ for all $k^* \leq k$. Then, by Markov's inequality,

$$3x^2 \max_{k^*(x) \leq k \leq n} \mathbf{P}\left(|S_k| \geq x\sqrt{n}/3\right) \leq \frac{3^4 \mathbf{E}|\,N(0,1)|}{x} = o_x(1)$$

for any $n \geq 1$. In particular,

$$3x^2 \limsup_{n \to \infty} \max_{k^*(x) \leq k \leq n} \mathbf{P}\left(|S_k| \geq x\sqrt{n}/3\right) = o_x(1)$$

Then, for $1 \leq k < k^*$ Chebyshev's inequality gives

$$3x^2 \limsup_{n \to \infty} \max_{0 \leq k < k^*} \mathbf{P}\left(|S_k| \geq x\sqrt{n}/3\right) \leq \limsup_{n \to \infty} \frac{3^3 k^*}{n} = 0$$

for any x. Altogether, this proves tightness by Lemma 3.8. $\square$

Remark. Since I had never seen it before, I will present Etemadi's inequality (a pretty tidy tool to have in your kit in my opinion). Let $(\xi_n)_{n=1}^\infty$ be a sequence of i.i.d. random variables, let $(S_n)_{n=0}^\infty$ be the partial sum of the first n ξ's, and let $t \geq 0$. Then, Etemadi's inequality states that

$$\mathbf{P}\left(\max_{1 \leq k \leq n} |S_k| \geq 3t\right) \leq 3 \max_{1 \leq k \leq n} \mathbf{P}(|S_k| \geq t).$$

With it, you can prove a weaker form of Kolmogorov's maximal inequality (one still strong enough to prove the strong law of large numbers though).

Remark. An entirely equivalent argument with A replacing 1 proves Donsker's Theorem on all compact sets of $[0,\infty)$, and hence proves that the result holds in the space $C[0,\infty)$ under the topology of uniform convergence on compact sets.

3.2 Skorokhod space

The question motivating the material we cover in this subsection is whether or not the convergence result of Donsker's Theorem can be extended to the far simpler sequence of functions for $0 \leq t \leq 1$,

$$J_n(t) = S_{\lfloor nt \rfloor} = \sum_{j=1}^{\lfloor nt \rfloor} \xi_j \quad \forall 0 \leq t \leq 1.$$

We should expect that J_n and W_n behave similarly, but in order to verify it we need to be able to talk about the convergence of random functions that are not continuous. Specifically, we study the set of all functions $f : [0, 1] \to \mathbb{R}$ such that

(i) For all $0 \leq t < 1$, $\lim_{s \downarrow t} f(s) = f(t)$.

(ii) For all $0 < t \leq 1$, $\lim_{s \uparrow t} f(s)$ exists.

Many know these functions by the name of càdlàg functions (I will write cadlag without the accents because I am lazy). We shall denote the space of all such functions by $D[0, 1]$. So, can we just argue our limits in the space $(D[0, 1], \| \cdot \|_\infty)$ and then move on to the next section?

The answer is no. The metric space $(D[0, 1], \| \cdot \|_\infty)$ is not very good for convergence. Specifically because functions with jump discontinuities that look almost identical can still be very far with respect to $\| \cdot \|_\infty$. This breaks properties like separability, which we leaned on for our weak convergence theory. For an example look no further than functions like $\mathbf{1}_{[a_n,1]}(x)$ for a sequence $(a_n)_{n=1}^\infty$ in $\mathbb{R}$ such that $a_n \to a$ as $n \to \infty$. If we are going to place a metric on functions with jump discontinuities such that the underlying metric space has nice characteristics, it should be the case that these functions to converge to $\mathbf{1}_{[a,1]}(x)$ as $n \to \infty$. So how do we remedy this situation? It is instructive to go back to the drawing board a bit and think about the properties that we want our metric d on $D[0, 1]$ to have. Here are some relatively reasonable desires:

(i) d encodes convergence of functions in $D[0, 1]$ in a way that allows convergence of sequences of functions like the example just mentioned above.

(ii) $(D[0, 1], d)$ is a complete separable metric space.

(iii) The restriction of d to the space $C[0, 1]$ is equivalent to $\| \cdot \|_\infty$.

(iv) If $\|x_n - x\|_\infty \to 0$ as $n \to \infty$ for $(x_n)_{n=0}^\infty$ a sequence in $D[0, 1]$, then $d(x_n, x) \to 0$ as $n \to \infty$.

(v) We can develop some theorems about d analogous to the Arzelà-Ascoli Theorem from the previous section.

The last point is probably the most confusing. The reason that we want this point to be true is that, without it, we would not have nice ways to prove tightness. Finding a metric that satisfies all of the above points seems like it may be a bit hard, but a fairly natural choice of d ends up giving us what we need. The idea is to view the domain $[0, 1]$ as being time moving forward, and to allow ourselves a very light amount of control over how fast time moves for each function.

Definition 3.10. *Let Λ be the set of all strictly increasing continuous functions from $[0, 1]$ onto itself. Then, we define the Skorokhod distance between two functions $x, y \in D[0, 1]$ to be given by*

$$d_S(x, y) = \inf_{\lambda \in \Lambda} \left(\max\{\|\lambda - I\|_\infty, \|x - y\lambda\|_\infty\}\right),$$

where $I : [0, 1] \to [0, 1]$ is the identity function, and $y\lambda(t) = y(\lambda(t))$ for all $0 \leq t \leq 1$. The pair $(D[0, 1], d_S)$ is called Skorokhod space.

A well known fact about cadlag functions is that they only have countably many jumps and only finitely many jumps above any height $\epsilon > 0$. This fact implies that cadlag functions are bounded, and hence the Skorokhod distance is always finite. We mention this just as a light sanity check for the definition.

For this subsection we are going to skip past the proofs and present the results that are true. Do not consider this a proper rigourous exploration of Skorokhod space, but instead as just a brief introduction to make the phrase less frightening when encountered in the wild. To see much full explanations, I would recommend looking at [Bil13, Ker22]. They were both excellent references for me when I was trying to undestand the concepts.

Most of the properties in (i)-(v) hold for d_S, but there is one small modification we need to make in order to have (ii) be true. As written, the sequence of indicators $\mathbf{1}_{[0,2^{-n})}(x)$ still does not converge as $n \to \infty$. You can check this by observing that any limit f of the sequence must have $f(x) = 0$ for all $x \in (0, 1]$, and then comparing this with the fact that for any $\lambda \in \Lambda$ there is some $x \in (0, 1])$ such that $\mathbf{1}_{[0,2^{-n})}(\lambda(x)) = 1$. The fix for this lack of completeness is to put a different condition on the time function λ. The second metric is given by

$$d_S^\circ(x, y) = \inf_{\lambda \in \Lambda} \left(\max\{\|\lambda\|^\circ, \|x - y\lambda\|_\infty\}\right),$$

where

$$\|\lambda\|^\circ = \sup_{0 \leq s < t \leq 1} \left| \log\left(\frac{\lambda(t) - \lambda(s)}{t - s}\right) \right|.$$

The two metrics d_S and d_S° are equivalent. That is a sequence in $D[0, 1]$ converges with respect to one if and only if it does with respect to the other (a stronger condition that is true is that they correspond to the same topology). Moreover, the metric space $(D[0, 1], d_S^\circ)$ is complete and separable (under the metric d_S we still have separability, just not completeness). For this second equivalent metric, all of the points (i)-(v) hold. The last of the five items that needs to be touched on is the last, (v). Rather than

Let $\Gamma(\delta) = \{(t_1, ..., t_k) : 0 \leq t_0, ..., t_k \leq 1,\ \min_{0\leq i\leq k}(t_i - t_{i-1}) > \delta\}$. Then, for $0 < \delta < 1$ we define

$$w'_x(\delta) = \inf_{(t_i)\in\Gamma(\delta)} \max_{1\leq i\leq k} \sup_{s,t\in[t_{i-1},t_i)} |x(s) - x(t)|.$$

Instead of just taking a supremum over the whole interval like with the modulus for $C[0, 1]$, we now allow ourselves to consider it over many small intervals. It is easy to see by taking the boundaries of the parition to be the points of the largest jump discontinuities that w'_x can be much smaller thatn w_x for functions in $D[0, 1]$. Our compactness characterization can be phrased in terms of this new w'_x function.

Theorem 3.11. *Let $S \subseteq D[0, 1]$. Then S is compact with respect to the Skorokhod topology if and only if*

(i) $\sup_{x\in S} \|x\|_\infty < \infty$.

(ii) $\lim_{\delta\to 0} \sup_{x\in S} w'_x(\delta) = 0$.

Now that we have successfully extended our ideas about $C[0, 1]$ to $D[0, 1]$, we need to say a couple things about measures on $D[0, 1]$. Specifically, we need to know about the relationship between a measure on $D[0, 1]$ and its finite dimensional marginals. The first thing to remark is that a coordinate projection π_t for some $0 < t < 1$ can is continuous at a function $x \in D[0, 1]$ if and only if the funciton is itself continuous at that point. This can be verified by appealing to the definition of the Skorokhod metric. This means that the question of measurabiltiy for the mappings is not even immediate. One can use the fact that the points of discontinuity for a function in $D[0, 1]$ form a Lebesgue measure zero set to argue that $\int_t^{t+\epsilon} x(s)\ ds$ is continuous for any $\epsilon > 0$ and hence measurable. Taking the limit as $\epsilon \to 0$ yields measurability for the coordinate projections. There main question is whether the coordinate projections form a separating class, i.e., whether measures with the same finite dimensional marginals are the same. This assertion is true, and leads us to a modified version of our characterization of weak convergence for $C[0, 1]$.

Theorem 3.12. *Let $(\mu_n)_{n=0}^\infty$ be a sequence of probability measures on $D[0, 1]$, and μ a measure on $D[0, 1]$. Let $T_\mu = \{0 \leq t \leq 1 : \mu(\{x : x(t) \neq \lim_{s\uparrow t} x(s)\}) = 0\}$. Then, if $\pi_{t_1,...,t_k}\mu_n \xrightarrow{\mathcal{L}} \pi_{t_1,...,t_k}\mu$ as $n \to \infty$ for all $t_1, ..., t_k \in T_\mu$ and $(\mu_n)_{n=0}^\infty$ is tight, then $\mu_n \xrightarrow{\mathcal{L}} \mu$ as $n \to \infty$.*

There is a lot more to say about the theory of weak convergence in $D[0, 1]$, but we have at least arrived at some analogous statements to those at the very beginning of the section. The finite dimensional marginal argument for our modified Donsker's Theorem carries over perfectly from the previous subsection. The tightness needs some more work that we do not have the time to get into, but to be able to use it later we conclude by presenting our second version of Donsker's Theorem.

Theorem 3.13. *Let $(\xi_n)_{n=0}^{\infty}$ be a sequence of i.i.d. random variables with mean 0 and variance 1. Then, $(J_n(t) : 0 \leq t \leq 1) \xrightarrow{\mathcal{L}} (B(t) : 0 \leq t \leq 1)$ as $n \to \infty$ in $(D[0,1], d_S^{\circ})$, where J_n is defined as at the beginning of this subsection and B is standard one dimensional Brownian motion.*

3.3 Convergence of the height process for Bienaymé forests

We now have everything we need to start exploring relationships between random combinatorial trees and random real trees. In this subsection, we start by showing that the height process of a critical Bienaymé tree converges to a Brownian excursion (see remark below). Let $(T_n)_{n=1}^{\infty}$ be a sequence of independent Bienaymé(μ) distributed random variables for some critical offspring distribution μ. Throughout the rest of this section we assume that all child distributions are critical. Let $X_i = |T_1| + ... + |T_i|$ for all $i \geq 1$. We define the height process of the forest by setting $H_k = h_{T_i}(k - X_{i-1})$ for all $X_{i-1} \leq k < X_i$ (recall that the height process of a tree $\mathbf{t}$ is defined on $0, ..., |\mathbf{t}| - 1$). Since the height process visits zero only once, the height process encodes the whole forest.

Before getting to the main theorem let's pause to address why the height function is the one we need to analyze. Our end goal is to prove the convergence of Bienaymé trees (specifically conditioned ones) to the Brownian CRT. To do this with Theorem 2.9, we need to show that the contour function of the tree converges to a Brownian excursion in distribution. We study the height function instead of the contour function is that the height function enjoys a nice connection with the DFQ process, which is distributed like a simple random walk for Bienaymé trees. Extending the result to include convergence of contour functions does not take much extra work. Of course, the desire to instead study the height function is what leads us to explore the convergence in Skorokhod space rather than $(C[0,1], \|\cdot\|_{\infty})$.

Remark. A Brownian excursion is, informally, a Brownian motion that is conditioned to be non-negative and takes the value 0 at time 1. This event of course has probability zero of occurring so we should be more careful than this. There are many legal ways to generate such stochastic processes, but one simple one goes as follows: Let $\tau_1, \tau_2 > 0$ be such that $B(\tau_1) = B(\tau_2) = 0$, $B(t) \geq 0$ for all $\tau_1 < t < \tau_2$ and $\tau_2 - \tau_1 \geq 1$ for some Brownian motion $(B(t) : t \geq 0)$. These times exist almost surely as Brownian motion is recurrent with expected return time to zero being unbounded. Then, set $e(t) = B((\tau_2 - \tau_1)t + \tau_1)/\sqrt{\tau_2 - \tau_1}$ for each $0 \leq t \leq 1$. This gives us a stochastic process with the correct characteristics.

Much of the work on combinatorial trees from Section 1 can be summarized with the following lemma.

Lemma 3.14. *For all $n \geq 0$, $H_n = |\{0 \leq k \leq n-1 : S_k = \inf_{k \leq j \leq n} S_j\}|$, where $(S_n)_{n=0}^{\infty}$ is a simple random walk with jump distribution ν defined by $\nu(k) = \mu(k+1)$ for all $k \geq -1$.*

Proof. Note that for $X_{i-1} \leq k < X_i$, the indices in $\{0 \leq k \leq n-1 : S_k = \inf_{k \leq j \leq n} S_j\}$ must be at least X_{i-1}. This is because each new tree is marked by a new global minimum in the random walk $(S_n)_{n=0}^{\infty}$. In particular, the kth tree ends where the random walk first visits the state $-k$. Thus, H_n coincides with h_{t_i} for $X_{i-1} \leq k < X_i$. From here, applying Theorems 1.6 and 1.10 complete the proof. □

Here is the main theorem.

Theorem 3.15. $(H_{\lfloor nt \rfloor}/\sqrt{n} : t \geq 0) \xrightarrow{\mathcal{L}} (2Z(t)/\sigma : t \geq 0)$ *as* $n \to \infty$*, where* σ^2 *is the variance of* μ*, and* $(Z(t) : t \geq 0)$ *is a reflected Brownian motion. The convergence occurs in* $D[0, \infty)$.

Remark. Reflected Brownian motion is $B(t) - \inf_{0 \leq s \leq t} B(s)$ for each $t \geq 0$, where $(B(t) : t \geq 0)$ is standard one dimensional Brownian motion. It has been study as far back as Lévy, and it is known to be distributed as $|B(t)|$.

Much of the heavy lifting in the proof of Theorem 3.15 is done by a couple of technical lemmas about random walks and a nice concentration inequality for the height process. We separate these pieces into their own pieces and then quickly explain why this completes the proof at the end. There exists proofs for the statement in full generality [Ald93], but they are not fully optimized to be able to present in a reasonable amount of time. For this proof we make one simplifying assumption that allows for the proving of the aforementioned concentration inequality we need. We assume that there is some $t > 0$ such that $\sum_{k \geq 0} \exp(tk)\mu(k) < \infty$, i.e., we assume that the moment generating function exists on some interval in the postive reals.

A few new pieces of notation need to be introduced before continuing. For the random walk defined in Lemma 3.14, $M_n := \sup_{0 \leq k \leq n} S_k$ and $I_n := \inf_{0 \leq k \leq n} S_k$. For all $n \geq 0$, we define the time reversed random walk starting from n by $\hat{S}_k^n := S_n - S_{n-k}$ for all $0 \leq k \leq n$. The duality principle for random walks asserts that $(\hat{S}_k^n : 0 \leq k \leq n) \stackrel{\mathcal{L}}{=} (S_k : 0 \leq k \leq n)$. For any sequence $x = (x_n)_{n=0}^{m}$ (m can be ∞), we define

$$\Phi_n(x) = \left| \left\{ 1 \leq k \leq n : x_k = \sup_{0 \leq j \leq k} x_j \right\} \right|.$$

Note that we do not count $k = 0$ in the size of the set. We can rewrite our expression for H_n in terms of our new notation.

Lemma 3.16. $H_n = \Phi_n(\hat{S}^n)$ *for all* $n \geq 0$.

Proof. Indeed,

$$\begin{aligned} S_k = \inf_{k \leq j \leq n} S_j &\iff S_n - S_k = S_n - \inf_{k \leq j \leq n} S_j \\ &\iff \hat{S}_{n-k}^n = \sup_{k \leq j \leq n} (S_n - S_{n-(n-j)}) \\ &\iff \hat{S}_{n-k}^n = \sup_{0 \leq j \leq n-k} \hat{S}_j^n. \end{aligned}$$

Thus, the cardinalities defining both functions (using the definition from Lemma 3.14) are the same. $\square$

Lemma 3.17. *Let* $(\tau_n)_{n=0}^{\infty}$ *be a sequence of stopping times defined inductively by setting* $\tau_0 = 0$ *and* $\tau_j = \inf\{n > \tau_{j-1} : S_n = M_n\}$ *for all* $j > 0$*. The sequence random variables* $(S_{\tau_j} - S_{\tau_{j-1}})_{j=1}^{\infty}$ *are i.i.d. with distribution given by*

$$\mathbf{P}(S_{\tau_1} - S_{\tau_0} = k) = \nu[k, \infty) = \mu[k+1, \infty)$$

for all $k \geq 0$.

Proof. The independence property is and immediate consequence of the Markov property. Let $R = \inf\{n \geq 1 : S_n = 0\}$ and let $k \in \mathbb{Z}$. Let $(\sigma_n)_{n=0}^{\infty}$ be the times at which the random walk is at either the state 0 or state k. The sequence $(S_{\sigma_n})_{n=0}^{\infty}$ is a symmetric Markov chain on the state space $\{0, k\}$. In particular, $\mathbf{E}T_{0,0} = 2$ as the stationary distribution is uniform. Hence, we expect to visit k once before returning to 0. Altogether, this shows that

$$\mathbf{E}\left[\sum_{n=0}^{R-1} \mathbf{1}_{\{S_n = k\}}\right] = 1. \tag{3}$$

Now, note that $\tau_1 \leq R$. If $S_1 > 0$, then $\tau_1 = 1$, and $R > 1$. If $S_1 < 0$, then τ_1 is the first time that the random walk is ≥ 0, which contains the event that the random walk returns to the origin. Moreover, since negative jumps of the walk are at most -1, the portion of the random walk on (τ_1, R) is all positive integers and the portion on $(1, \tau_1)$ is all negative integers. In particular, if the value k defining R is taken to be nonpositive, then by (3),

$$\mathbf{E}\left[\sum_{n=0}^{\tau_1 - 1} f(S_n)\right] = \sum_{i=0}^{\infty} f(-i)\mathbf{E}\left[\sum_{n=0}^{R-1} \mathbf{1}_{\{S_n = -i\}}\right] = \sum_{i=0}^{\infty} f(-i) \tag{4}$$

for any function $f : \mathbb{Z} \to \mathbb{Z}_{\geq}$. Continuing,

$$\begin{aligned}
\mathbf{E}[f(S_{\tau_1})] &= \sum_{n=0}^{\infty} \mathbf{E}\left[f(S_{n+1})\mathbf{1}_{\{n<\tau_1\}\cap\{S_{n+1}\geq 0\}}\right] \\
&= \sum_{n=0}^{\infty}\sum_{j=0}^{\infty} \mathbf{E}\left[f(S_n + j)\nu(j)\mathbf{1}_{\{n<\tau_1\}\cap\{S_n+j\geq 0\}}\right] \\
&= \sum_{j=0}^{\infty} \nu(j)\mathbf{E}\left[\sum_{n=0}^{\infty} f(S_n + j)\mathbf{1}_{\{n<\tau_1\}\cap\{S_n+j\geq 0\}}\right] \\
&= \sum_{j=0}^{\infty} \nu(j)\mathbf{E}\left[\sum_{n=0}^{\tau_1 - 1} f(S_n + j)\mathbf{1}_{\{S_n+j\geq 0\}}\right]
\end{aligned}$$

$$= \sum_{j=0}^{\infty} \sum_{i=0}^{\infty} \nu(j) f(j-i) \mathbf{1}_{\{j \geq i\}} \qquad \text{(by (4))}$$

$$= \sum_{m=0}^{\infty} \sum_{\ell=m}^{\infty} f(m) \nu(\ell).$$

From here, just take for all $k \in \mathbb{Z}$, $f(x) = \mathbf{1}_{\{x=k\}}$ to obtain the desired result. □

With this, we can prove a key part of the proof of Theorem 3.15.

Lemma 3.18.

$$\frac{H_n}{S_n - I_n} \xrightarrow{\mathbb{P}} \frac{2}{\sigma^2}$$

as $n \to \infty$.

Proof. Let the sequence $(\tau_n)_{n=0}^{\infty}$ be defined as above. As ν has mean zero,

$$\mathbf{E}[S_{\tau_1}] = \sum_{k=0}^{\infty} k\nu[k, \infty) = \sum_{j=0}^{\infty} \frac{j(j+1)}{2} \nu(j) = \sigma^2/2.$$

Moreover,

$$M_n = \sum_{k:\tau_k \leq n} (S_{\tau_k} - S_{\tau_{k-1}}) = \sum_{k=1}^{\Phi_n(S)} (S_{\tau_k} - S_{\tau_{k-1}}).$$

By Lemma 3.17 and the law of large numbers, $M_n/\Phi_n(S) \xrightarrow{\text{a.s.}} \sigma^2/2$ as $n \to \infty$ ($\Phi_n(S) \to \infty$ almost surely as $n \to \infty$ by null recurrence). Using Lemma 3.16 and the duality principle, we have that $(M_n, \Phi_n(S)) \stackrel{\mathcal{L}}{=} (S_n - I_n, H_n)$ for all $n \geq 0$. Hence,

$$\frac{S_n - I_n}{H_n} \xrightarrow{\mathbb{P}} \frac{\sigma^2}{2}$$

as $n \to \infty$. □

Now we turn our attention to the issue of concentration. In the proof of Theorem 3.15 we use a stronger result than just the law of large numbers convergence from the previous proof. Given the previous two results, the proof is not too different from that for most standard concentration inequalities in probabilistic combinatorics. A full proof can be found in [LG05], we shall just record the result and move on.

Lemma 3.19. *For any* $\epsilon \in (0, 1/4)$ *there exists a* $\delta > 0$ *and an* $N \geq 1$ *such that for all* $n \geq N$ *and all* $0 \leq j \leq n$,

$$\mathbf{P}\left(\left| M_j - \frac{\sigma^2}{2} \Phi_j(S) \right| \geq n^{1/4+\epsilon} \right) \leq e^{-n^{\delta}}.$$

We are now ready to prove Theorem 3.15.

Theorem. $(H_{\lfloor nt \rfloor}/\sqrt{n} : t \geq 0) \xrightarrow{\mathcal{L}} (2Z(t)/\sigma : t \geq 0)$ *as* $n \to \infty$*, where* σ^2 *is the variance of* μ*, and* $(Z(t) : t \geq 0)$ *is a reflected Brownian motion. The convergence occurs in* $D[0, \infty)$ *with its associated metric.*

Proof. Most of the tough computations were done in the above lemmas. We just need to carefully go through and check that all of the convergences line up in the right way.
Step 1: (The function $\varphi : D[0, A] \to D[0, A]$ defined by $\varphi(f)(t) = \sup_{0 \leq s \leq t} f(s)$ is continuous with respect to the Skorokhod topology) Suppose that x, y are such that $d(x, y) < \delta$ and without loss of generality assume that there is no dilation (of course, we could just redefine y to be λy). Let $t \in [0, A]$ and suppose without loss of generality that $\sup_{0 \leq s \leq t} x(s) \geq \sup_{0 \leq s \leq t} y(s)$. Let $(s_k)_{k=1}^{\infty}$ be such that $x(s_k) \to \sup_{0 \leq s \leq t} x(s)$. We have for large k that $\delta \leq y(s_k) \leq x(s_k)$. By compactness, we may take some subsequence $(s_{k_m})_{m=1}^{\infty}$ such that $y(s_{k_m}) \to \alpha^*$ for some α^*. Then, it must hold that $\sup_{0 \leq s \leq t} x(s) - \delta \leq \alpha^* \leq \sup_{0 \leq s \leq t} y(s) \leq \sup_{0 \leq s \leq t} x(s)$. Since t was chosen arbitrarily the result follows.

Step 1, Donsker's Theorem, and the continuous mapping theorem combine to give that

$$\left(\frac{1}{\sqrt{n}} (S_{\lfloor nt \rfloor} - I_{\lfloor nt \rfloor}) : t \geq 0 \right) \xrightarrow{\mathcal{L}} \left(\sigma(B(t) - \inf_{0 \leq s \leq t} B(s)) : t \geq 0 \right)$$

as $n \to \infty$ in $D[0, \infty)$ (recall that convergence in $D[0, \infty)$ is equivalent to convergence in $D[0, A]$ for all values of A).

Step 2: (Turning $S - I$ into H) Recall from the proof of Lemma 3.18 that $(S_n - I_n, H_n) \stackrel{\mathcal{L}}{=} (M_n, \Phi_n(S))$. Thus, Lemma 3.19 implies that for all $0 \leq j \leq n$ for n large that

$$\mathbf{P} \left(\left| S_j - I_j - \frac{\sigma^2}{2} H_j \right| > n^{3/8} \right) \leq e^{-n^{\epsilon'}}$$

for some $\epsilon' > 0$. An application of the union bound gives

$$\mathbf{P} \left(\sup_{0 \leq j \leq n} \left| S_j - I_j - \frac{\sigma^2}{2} H_j \right| > n^{3/8} \right) \leq n e^{-n^{\epsilon'}}.$$

We can easily extend the event to the continuous height function on the interval $[0, A]$,

$$\mathbf{P} \left(\sup_{0 \leq t \leq A} \left| S_{\lfloor nt \rfloor} - I_{\lfloor nt \rfloor} - \frac{\sigma^2}{2} H_{\lfloor nt \rfloor} \right| > (An)^{3/8} \right) \leq A n e^{-(An)^{\epsilon'}}.$$

Summing and applying the Borel-Cantelli lemma we get that

$$\sup_{0 \leq t \leq A} \left| \frac{S_{\lfloor nt \rfloor} - I_{\lfloor nt \rfloor}}{\sqrt{n}} - \frac{H_{\lfloor nt \rfloor}}{\sqrt{n}} \right| \xrightarrow{\text{a.s.}} 0$$

as $n \to \infty$. Combining this with the conclusion after step 1 yields the final result. $\square$

3.4 Convergence of the contour process

Towards the goal of proving convergence in the Gromov-Hausdorff topology, we would also like to say something about the convergence of contour functions for trees (and forests). Luckily this follows quite easily from the convergence for the height process. In this subsection, we give a contour function analogue for Theorem 3.15.

If we want to make a contour process for a sequence of independent Bienaymé(μ) trees $(T_n)_{n=1}^{\infty}$ then we need to deal with the fact that the contour function for the tree $\{\emptyset\}$ is trivial. Recall that the contour function $\gamma_{\mathbf{t}}$ for a tree $\mathbf{t}$ is defined on the interval $[0, 2(|\mathbf{t}|-1)]$. We define a new contour function $\gamma'_{\mathbf{t}}$ by $\gamma'_{\mathbf{t}}(t) = \gamma(t)\mathbf{1}_{\{t\in[0,2(|\mathbf{t}|-1)]\}}$. We define the contour process $(\Gamma(t) : t \geq 0)$ by concatenating the functions $(\gamma'_{T_n})_{n=1}^{\infty}$. For all $n \geq 0$ define $J_n = 2n - H_n + I_n$, where we recall that $I_n = \sup_{0\leq k\leq n} S_k$.

Lemma 3.20. *Let $(T_n)_{n=1}^{\infty}$ be a sequence of independent Bienaymé(μ) trees with $(U_n)_{n=0}^{\infty}$ being the vertices written in lexicographical order (that is, the ordering obtained from making the root of T_{n+1} larger than every vertex of T_n for all $n \geq 1$ and ordering individual trees with the standard lexicographical order). Then, over the interval $[J_n, J_{n+1}]$ the contour process goes from the height of U_n to the height of U_{n+1}. Moreover,*

$$\sup_{t\in[J_n,J_{n+1}]} |\Gamma(t) - H_n| \leq |H_{n+1} - H_n| + 1.$$

Proof. There are three possible cases (proof by look at Figure 5):

(i) U_{n+1} is a child of U_n;

(ii) U_{n+1} is a child of an ancestor of U_n;

(iii) U_{n+1} is the root of the next tree in the sequence.

It is pretty straightforward to verify both the first and the second statements by induction by splitting them into these cases. □

Theorem 3.21. *If $(\Gamma(t) : t \geq 0)$ is the contour process for a sequence of Bienaymé(μ) random forests, then*

$$\left(\frac{1}{\sqrt{n}}\Gamma(2nt) : t \geq 0\right) \xrightarrow{\mathcal{L}} \left(\frac{2}{\sigma}Z(t) : t \geq 0\right)$$

in $D[0,\infty)$ as $n \to \infty$, where $(Z(t) : t \geq 0)$ is reflected Brownian motion.

Proof. Let $A > 0$. Let $\varphi : [0,\infty) \to \mathbb{N}$ be a random function defined by $\varphi(t) = \sum_{n=0}^{\infty} n\mathbf{1}_{\{t\in[J_n,J_{n+1})\}}$. By Lemma 3.20 and Theorem 3.15,

$$\sup_{t\leq A}\left|\frac{1}{\sqrt{n}}\Gamma(2nt) - \frac{1}{\sqrt{n}}H_{\varphi(2nt)}\right| \leq \frac{1}{\sqrt{n}} + \frac{1}{\sqrt{n}}\sup_{t\leq A}|H_{\lfloor nt\rfloor+1} - H_{\lfloor nt\rfloor}| \xrightarrow{\mathbb{P}} 0$$

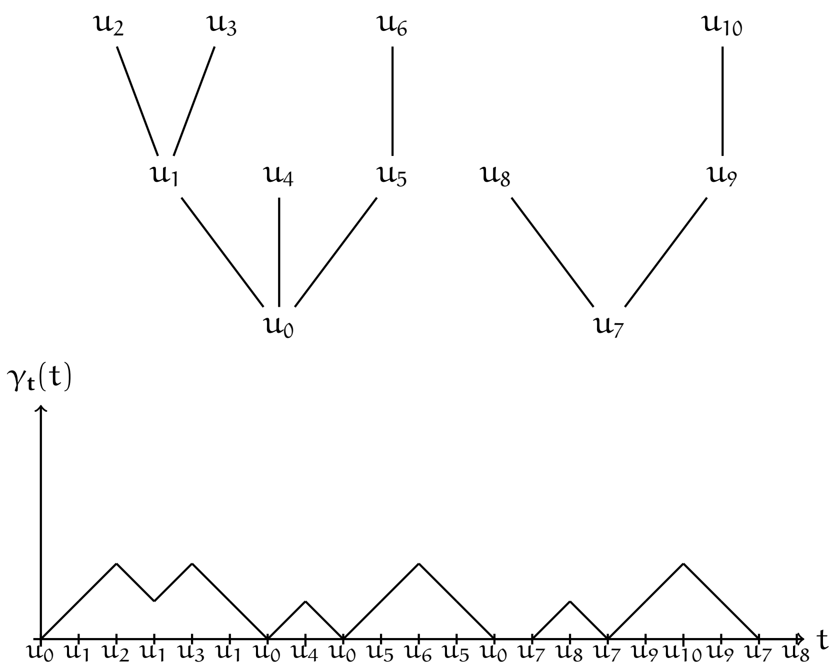

Figure 5: The first two trees in a realization of a Bienaymé forest along with the contour process $\Gamma(t)$ for the first two trees.

as $n \to \infty$. From the definition of the intervals $(J_n)_{n=0}^{\infty}$ we have that, for any $m \geq 0$

$$\sup_{0\leq t\leq m}\left|\varphi(t)-\frac{t}{2}\right| \leq \sup_{n:J_n\leq m}\sup_{t\in[J_n,J_{n+1})}\left|n-\frac{t}{2}\right| \leq \sup_{n:J_{n-1}\leq m}\left|\frac{J_n}{2}-n\right|+1 \leq \sup_{n:J_{n-1}\leq m}\frac{H_n+|I_n|}{2}+1.$$

It is clear that $\{n : J_{n-1} \leq m\} \subseteq [m+1]$, and so

$$\sup_{0\leq t\leq m}\left|\varphi(t)-\frac{t}{2}\right| \leq \sup_{n\leq m+1}\frac{H_n}{2}+\frac{|I_m|}{2}+1.$$

Replacing m with $2An$ we get,

$$\frac{1}{n}\sup_{0\leq t\leq A}|\varphi(2nt)-nt| \leq \frac{1}{n}\sup_{0\leq k\leq 2An+1}H_k+\frac{1}{n}|I_{2An+1}|+\frac{1}{n}\xrightarrow{\mathbb{P}} 0$$

as $n \to \infty$. Combining the two inequalities and applying Theorem 3.15 one more time we arrive at the final result. □

3.5 Aldous' Theorem

We are ready to turn our attention to combinatorial trees again and prove our first scaling limit theorem for random trees. Specifically, we identify a universal limit for conditioned Bienaymeé trees. The universal limit is known as the Brownian continuum random tree.

Definition 3.22. *Let $(e(t) : t \in [0,1])$ be a Brownian excursion. Extend the function to $[0,\infty)$ by defining $e(t) = 0$ for $t > 1$. The random metric space $\mathbf{T}_e$ is called the Brownian continuum random tree (CRT).*

We shall learn about the CRT as we continue to develop the theory of scaling limits (specifically, Section 4 offers a lot of insight into the structure of the tree), though for the moment it's, main importance is that it is the limit in the following theorem.

Theorem 3.23 (Aldous' Theorem). *Let $\mathbf{T}_n \stackrel{\mathcal{L}}{=} \text{Bienaymé}(n,\mu)$ be a critical Bienaymé tree considered as a real tree with length one edge lengths (take the tree encoded by the contour function). If μ has finite variance σ^2, then $\frac{1}{\sigma\sqrt{2n}}\mathbf{T}_n \xrightarrow{\mathcal{L}} \mathbf{T}_e$ as $n \to \infty$ in the space $(\mathbb{T}, d_{GH})$.*

As one can imagine from the work done above, the convergence is essentially a corollary of a functional convergence result for the height/contour functions.

Theorem 3.24. *Let $T_n \stackrel{\mathcal{L}}{=} \text{Bienaymé}(n,\mu)$ be a non-trivial critical Bienaymé tree, and let σ^2 be the variance of μ. Let $(H_k^{(n)})_{k=1}^n$ be the height process for T_n for each $n \geq 1$. Then,*

$$\left(\frac{1}{\sqrt{n}} H^{(n)}_{\lfloor nt \rfloor} : 0 \leq t \leq 1\right) \xrightarrow{\mathcal{L}} \left(\frac{2}{\sigma} e(t) : 0 \leq t \leq 1\right)$$

as $n \to \infty$ in $D[0,1]$. $(e(t) : 0 \leq t \leq 1)$ is a normalized length 1 Brownian excursion.

The proof builds on the work done on Theorem 3.15, though some additional effort is needed to address the fact that the trees have a fixed size. This change removes the independence between each jump in the walk and breaks the ability to apply Donsker's Theorem. Because of this, we need a version of Donsker's Theorem for discrete excursions.

Lemma 3.25. *Let $(\xi_k)_{k=1}^\infty$ be a sequence of i.i.d. random variables with mean 0 and variance 1 and let $S_k = \sum_{i=1}^k \xi_i$ for all $k \geq 0$. Let $\tau = \inf\{k \geq 1 : S_k \leq 0\}$. Let $(S_k^*)_{k=0}^\infty$ be distributed like S_k under $\mathbf{P}(\cdot|\tau = n)$, i.e., $\mathbf{P}(S_k^* = j) = \mathbf{P}(S_n = j|\tau = n)$ for all $k \geq 0$.*

$$\left(\frac{1}{\sqrt{n}} S^*_{\lfloor nt \rfloor} : 0 \leq t \leq 1\right) \xrightarrow{\mathcal{L}} (e(t) : 0 \leq t \leq 1)$$

as $n \to \infty$ in $D[0,1]$.

The proof of this lemma follows a similar structure to the proof of the original, and was developed over many papers in the 1970's [Bel72, Kai75, Kai76]. If I have time later I might try to fill this proof in, but for now I'm going to skip past it.

Proof of Theorem 3.23. We shall deal only with the convergence of the height process for the trees $(T_n)_{n=1}^\infty$, noting that converting the result to be about the contour function follows the exact same structure as the conversion of Theorem 3.15 provided in

Theorem 3.21. Let $T \overset{\mathcal{L}}{=} \text{Bienaymé}(\mu)$ be an unconditioned tree and let $(S_n)_{n=0}^{\infty}$ and $(H_n)_{n=0}^{|T|-1}$ be its corresponding DFQ process and height process. From the local limit theorem for simple random walks we have that

$$\lim_{n\to\infty} \sup_x \left| \sqrt{2\pi n}\sigma \mathbf{P}(S_n = x) - e^{-\frac{x^2}{2n\sigma^2}} \right|.$$

Then, using the cycle lemma for simple random walks,

$$\mathbf{P}(|T| = n) = \mathbf{P}(S_0 \geq 0, ..., S_{n-1} \geq 0, S_n = -1) = \frac{1}{n}\mathbf{P}(S_n = -1) \sim \frac{1}{\sigma\sqrt{2\pi n^3}}.$$

In proving Theorem 3.15, we proved that

$$\mathbf{P}\left(\sup_{0\leq t\leq 1} \left| \frac{H_{\lfloor nt \rfloor}}{\sqrt{n}} - \frac{2(S_{\lfloor nt \rfloor} - I_{\lfloor nt \rfloor})}{\sigma^2\sqrt{n}} \right| > n^{-1/8} \right) \leq n e^{-n^{\epsilon}}$$

for some $\epsilon > 0$. As $\mathbf{P}(|T| = n)$ is polynomial in n we can condition without changing much,

$$\mathbf{P}\left(\sup_{0\leq t\leq 1} \left| \frac{H_{\lfloor nt \rfloor}}{\sqrt{n}} - \frac{2(S_{\lfloor nt \rfloor} - I_{\lfloor nt \rfloor})}{\sigma^2\sqrt{n}} \right| > n^{-1/8} \;\middle|\; |T| = n \right) \leq O\left(n^{5/2} e^{-n^{\epsilon}}\right).$$

Recalling the continuity of the supremum and infimum with respect to the Skorokhod topology and applying Lemma 3.25 we get that $(H^{(n)}_{\lfloor nt \rfloor} : 0 \leq t \leq 1) \xrightarrow{\mathcal{L}} (\frac{2}{\sigma^2} e(t) : 0 \leq t \leq 1)$ in the space $D[0,1]$, where we are defining $(H^{(n)}_n)_{n=0}^{n-1}$ to be the height process for T_n. □

A remark on the height of the Brownian CRT

A nice corollary of Aldous' Theorem is that the height of critical Bienaymé trees scaled by $1/\sqrt{n}$ converges to the height of the Brownian CRT. This naturally leads one to wonder what the height of the Brownian CRT is. Recall that the root of $\mathbf{T}_e$ is the equivalence class $[0]_{R_e}$. Since $e(0) = 0$, this implies that $\text{ht}(\mathbf{T}_e) = \sup_{0\leq t\leq 1} e(t)$. That is a pretty clean description, but studying $\sup_{0\leq t\leq 1} e(t)$ is far from an easy job. For example, the diameter (which is closely related to the height), has a probability density given by

$$f(y) = \frac{\sqrt{2\pi}}{3} \sum_{n\geq 1} \left(\frac{64}{y^2}\left(4b_{n,y}^4 - 36b_{n,y}^3 + 75b_{n,y}^2 - 30b_{n,y}\right) + \frac{16}{y^2}\left(2b_{n,y}^3 - 5b_{n,y}^2\right) \right) e^{-b_{n,y}},$$

where $b_{n,y} = \frac{8\pi^2 n^2}{y^2}$ [Sze06, Wan15]. It can be proven either by using the Brownian CRT's close relationship with combinatorial trees or via direct analysis of the the supremeum of Brownian excursions.

4 The line-breaking construction of the CRT

In this section we cover an alternative perspective from which one can prove Aldous' Theorem in the case that the tree is distributed uniformly over all trees of size n. Rather than using functional encodings, this proof is algorithmic, providing a more tangible way to both generate and study the Brownian CRT. The result is stronger than Aldous' Theorem, as the convergence is with respect to a generalization of the Gromov-Hausdorff distance from Section 2. In this generalization we equip our trees with measures, which give further information about the structure of the trees. Most of the material from this section was initially explored in [Ald90, Ald91].

4.1 The Gromov-Hausdorff-Prokhorov distance

In this section we prove convergence results with respect to a stricter topology than the one induced from the Gromov-Hausdorff distance. Specifically, it is a generalization of the distance to the space of measured metric spaces. Essentially, it just meant to encode both weak convergence of measure and metric space. For the convergence of measure part we need to introduce the Prokhorov metric (are you tired of new metrics yet?).

Definition 4.1. *For two measures μ and ν defined on a metric space (X, d), we define the Prokhorov distance between them to be*

$$d_P = \inf\{\epsilon > 0 : \mu(A) \leq \nu(A^\epsilon) + \epsilon, \nu(A) \leq \mu(A^\epsilon) + \epsilon \; \forall A \in \mathcal{B}(X)\},$$

where $\mathcal{B}(X)$ are the Borel-measurable sets and $S^\epsilon = \{x \in X : d(x, S) \leq \epsilon\}$.

The next lemma explains why we introduced another metric. Using the Portmanteau theorem it is not too hard to prove (see e.g., [Bil13]).

Lemma 4.2. *Let $M(X)$ be the space of finite measures on some complete separable metric space (X, d). Let $(\mu_n)_{n=1}^{\infty}$ be a sequence in $M(X)$ and let $\mu \in M(X)$. Then $\mu_n \xrightarrow{\mathcal{L}} \mu$ if and only if $d_P(\mu_n, \mu) \to 0$ as $n \to \infty$, i.e., d_P is a metrization of the topology of weak convergence.*

Now we can formally introduce the idea of Gromov-Hausdorff-Prokhorov convergence (GHP). The Prokhorov distance is needed in order to have a well defined metric that encapsulates what we want.

Definition 4.3. *Let* (X, d_X, μ) *and* (Y, d_Y, ν) *be compact measured metric spaces with roots* ρ_X *and* ρ_Y*, where* $\mu \in M(X)$ *and* $\nu \in M(Y)$*. We define the Gromov-Hausdorff-Prokhorov distance between* X *and* Y *as*

$$d_{GHP}(X, Y) = \inf_{\varphi_X, \varphi_Y} \left(\max\{d_H(\varphi_X(X), \varphi_Y(Y)), d^*(\varphi_X(\rho_X), \varphi_Y(\rho_Y)), d_P(\varphi_X\mu, \varphi_Y\nu)\} \right),$$

where the infimum is taken over all isometric embeddings φ_X *and* φ_Y *into a metric space* (X^*, d^*).

The reason that convergence in the GHP topology is preferred in many cases is quite simple - measures on metric spaces can tell us information about the space if the measure is a function of the structure. Specifically, we care a lot about the uniform probability measure on a compact space. We are going to try to keep rigorous discussion of the GHP topology to a minimum and refer the reader to [ADH13] if they are interested in the finer details. We'll record below that the metric induces some good metric spaces as was done with the GH metric. There is an equivalent definition of GHP distance that induces the same topology and uses the notion of correspondences instead of isometries that one can find in [ABBGM17]. For the next theorem, we say that two measured compact metric spaces (X, d, μ) and (Y, d, ν) are called *equivalent* if there is an isometry $\varphi : X \to Y$ such that $\varphi\mu = \nu$.

Theorem 4.4. *If* $\mathbb{K}_M$ *is the space of equivalence classes of measure compact metric spaces, then* $(\mathbb{K}_M, d_{GHP})$ *is a complete separable metric space as well.* $\mathbb{T}_M$*, the space of equivalence classes of measured real trees, is a closed set in* $\mathbb{K}_M$

4.2 Growing random trees (the set up)

Consider the following combinatorial tree growth algorithm for a fixed $n \geq 2$ and vertex set V_n labelled in $[n]$ and edge set E_n:

(i) For $2 \leq i \leq n$, let $V_i = \min\{U_i, i-1\}$, where $U_i \overset{\mathcal{L}}{=} \text{Unif}(\{1, ..., n\})$. Add the edge $\{i, V_i\}$ to E_n.

(ii) Remove the labels from V_n and then give them a new label distributed uniformly over all bijections $V_n \to [n]$.

(iii) output $T_n = (V_n, E_n)$.

As the technique deviates a bit from those in this course, we won't prove the following theorem that can be found in [Ald90]. The proof relies on connections between uniform spanning trees of graphs and random walks on them.

Theorem 4.5. *Let* T_n *be the output from the above algorithm. Then,* T_n *is distributed like a uniform lablled tree on* n *vertices.*

Now, consider the following real tree growth "algorithm" that is done within the space ℓ_1 of absolutely summable sequences with the metric $d(x, y) = \sum_{n=1}^{\infty} |x_n - y_n|$. Let $C_0 = L_0 = 0$, define $(C_n)_{n=1}^{\infty}$ to be the arrival times of a non-homogenous poisson process with rate $r(t) = t$, and define $(L_n)_{n=1}^{\infty}$ by $L_n = C_n \cdot \xi_n$, where ξ_n is a Unif[0, 1] random variable. Then define $f(0) = 0$ and

$$f(x) = \sum_{n=1}^{\infty} (f(L_{i-1}) + (x - C_{i-1})z_i)\mathbf{1}_{\{C_{i-1} < x \le C_i\}},$$

where z_i is the basis vector in ℓ_1 in the canonical basis for the space that has the ith entry non-zero. Then, we set $\mathcal{T}_x = f([0, x])$ for all $x \ge 0$, $\mathcal{T}_\infty = f([0, \infty))$, and $\mathcal{T} = \overline{\mathcal{T}_\infty}$ (the closure). These two procedures are the main stars of this section. Let's just present all the results other in a row.

Theorem 4.6. *Let $(\mu_t : t \ge 0)$ be defined by setting $\mu_t = f\lambda_{[0,t]}$ for all $t \ge 0$ ($\lambda_{[0,t]}$ is the uniform probability measure on $[0, t]$ and f is the function definining the continuum algorithm). Then, almost surely, there is some measure $\mu_\mathcal{T}$ such that $\mu_t \xrightarrow{\mathcal{L}} \mu_\mathcal{T}$ as $t \to \infty$. Moreover, $\mu_\mathcal{T}$ is supported on $\mathcal{T}$ and satisfies $\mu_\mathcal{T}(\mathcal{T} \setminus \mathcal{T}_\infty) = 1$ almost surely.*

Theorem 4.7. *Let $(T_n)_{n=2}^{\infty}$ be defined as in the combinatorial algorithm (viewed as a real tree), let $\mathcal{T}$ be defined as the continuum algorithm, and let μ_n be the uniform measure on vertices in T_n (viewed as a measure on the real tree version of T_n). Then, $(T_n, \mu_n) \xrightarrow{\mathcal{L}} (\mathcal{T}, \mu_\mathcal{T})$ as $n \to \infty$ with respect to the GHP topology. In particular, $\mathcal{T}$ is distributed like the Brownian CRT.*

Corollary 4.8. *The Brownian CRT has only vertices of degree one, two, and three almost surely. The uniform measure on it is supported almost surely on the set of leaves.*

The proof strongly utilizes the nice properties of Poisson procceses to learn about properties of the tree $\mathcal{T}$. In the previous section, continuous functions were our window into the the world of real trees. Here, Poisson processes can be seen as playing a similar role.

An aside on non-homogeneous Poisson processes

Crucial to the continuum algorithm is a non-homogeneous Poisson process on $[0, \infty)$. It is worthwhile to review some basic facts about them. Let $r : [0, \infty) \to [0, \infty)$ be Lebesgue measurable and define

$$m(t) = \int_0^t r(s)\ ds.$$

for all $t \ge 0$. Similarly, we define for a measurable set A,

$$m(A) = \int_A r(s) ds.$$

Let $(C_k)_{k=1}^{\infty}$ be a collection of random points in $[0, \infty)$, and let $N_t = |\{k : C_k \le t\}|$. Define $N(A)$ for a measurable set analogously. We say that the collection is a *Poisson process with rate function* r if, for any collection of disjoint measurable sets $(A_k)_{k=1}^{\infty}$, the collection $(N(A_k))_{k=1}^{\infty}$ consists of independent Poi$(m(A_k))$ random variables. One can argue that such processes exist by studying regular Poisson processes under a time change. One big fact that we need for the proceeding proof is about how many points can arrive in small intervals of time.

Theorem 4.9. *Let $(C_k)_{k=1}^{\infty}$ be a Poisson process with rate function $r(t)$ with counting process N. Then, for all $t \ge 0$,*

(i) $\lim_{h\downarrow 0} \frac{\mathbf{P}(N[t,t+h)=1)}{h} = r(t)$, *and*

(ii) $\lim_{h\downarrow 0} \frac{\mathbf{P}(N[t,t+h)>1)}{h} = 0$.

The other topic that we need to be familiar with is the inter-arrival times of the Poisson process. There is a general formula for the pdf of the random variable C_k given by

$$f_k(t) = \frac{(m(t))^{k-1}}{(k-1)!} r(t) e^{-m(t)}.$$

In particular, $\mathbf{P}(C_1 > t) = e^{m(t)}$. Any sequence of random variables that satisfies these properties is a non-homogeneous Poisson process with rate $r(t)$.

4.3 THE FRACTAL DIMENSION OF THE TREE $\mathcal{T}$

The first thing to cover is the question of compactness for the tree $\mathcal{T}$. For the sake of brevity, we just outline the stucture of the proof, leaving out almost all of the details. The approach towards proving $\mathcal{T}$ is compact also happens to proves that the Minkowski dimension of the tree is 2. Here we define the *Minkowski dimension* of a set S to be

$$\lim_{\epsilon\downarrow 0} \frac{\log N(S, \epsilon)}{\log(\epsilon^{-1})},$$

where $N(S, \epsilon)$ is the number of balls of radius ϵ needed to cover S. In particular, $\mathcal{T}$ is compact. For the rest of the section we shall concern ourselves mostly with the function

$$D(s, t) = \inf_{0 \le r \le s} \|f(t) - f(r)\|$$

whre $0 < s < t$. $D(s, t)$ is the distane between the point $f(t)$ and the tree $\mathcal{T}_s$. Some sort of summable bound on the probability that $\sup_{t \ge s} D(s, t)$ is large should be enough to assert compactness. We first turn our attention to a couple computational lemmas that give us exactly such a bound.

Lemma 4.10. *For $0 < s < t$, $D(s, t)$ is stochastically dominated by an* Exp(s) *random variable. Moreover, there is a random function $(X(s) : s > 0)$ such that $D(s, t) \xrightarrow{\mathcal{L}} X(s)$ as $t \to \infty$ and $sX(s) \xrightarrow{\mathcal{L}}$* Exp$(1)$ *as $s \to \infty$.*

Proof. We are just going to cover the second point. Define $g(t) = \mathbf{E}[\exp(\theta D(s,t))]$ (θ and s are taken to be fixed). I'm going to do a mildly illegal move and assume that, if an arrival occurs in the interval $[t, t+h)$, it occurs immediately at t. This only gives a bound of course, but the argument in full generality is super messy and I think this computation at least gives the intuition behind the answer as good as doing the whole thing. Of course, to make everything properly rigourous, we would need to consider the arrival time over the set $[0, h]$.

Set $p_0 = \mathbf{P}(N[t, t+h) = 0)$ and $p_1 = \mathbf{P}(N[t, t+h) = 1)$ Using Theorem 4.9 and the independence properties of Poisson processes we get, for a uniform random variable U,

$$\begin{aligned}\frac{1}{h}(g(t+h) - g(t)) &\approx \mathbf{E}\left[e^{\theta(D(s,t\cdot U)+h)}\right]\frac{p_1}{h} + \mathbf{E}\left[e^{\theta(D(s,t)+h)}\right]\frac{p_0}{h} - \frac{g(t)}{h} + o(1)\\ &\approx \frac{p_1 e^{\theta h}\int_0^t g(u)\,du}{th} + \frac{p_0 e^{\theta h} g(t)}{h} - \frac{g(t)}{h} + o(1)\\ &\approx \frac{p_1 e^{\theta h}\int_0^t g(u)\,du}{th} + \frac{(1-p_1)(1+\theta h)g(t) - g(t)}{h} + o(1).\end{aligned}$$

If we take a limit as $h \to 0$ and replace approximations with equality we get

$$g'(t) = \theta g(t) - t g(t) + \int_0^t g(u)\,du.$$

Moreover, we have boundary conditions, $g(u) = 1$ for all $0 \le u \le s$ and $\lim_{u\downarrow s} g'(u) = \theta$. The function

$$g(t) = 1 + \frac{\theta}{\exp(\theta t - t^2/2)}\int_s^t \exp(\theta u - u^2/2)\,du$$

solves this system. One can observe the desired converges from just taking limits here and using the equivalence between convergence of the mgf and weak convergence. $\square$

From playing around with the stochastic domination in the above result we can derive the following maximal inequality.

Lemma 4.11. *With* $D(s,t)$ *define as above we have,*

$$\mathbf{P}\left(\sup_{e^j \le t \le e^{j+1}} D(e^j, t) \ge 3je^{-j}\right) \le e^{2-j}.$$

Slapping the Borel-Cantelli lemma on this bound and doing some algebra gives that the covering number $N(\mathcal{T}, \delta)$ is bounded above by $(13e\delta^{-1}\log(\delta^{-1}))^2$ almost surely for sufficiently small $\delta > 0$. This proves compactness and also proves that the Minkowski dimension is at most two. Let $A(t, \delta)$ be the events that no arrival occurs in the Poisson process in the interval $[t, t+2\delta]$. One can analyze the collection of points $M(\delta) = \{f(2\delta n) : n \geq 1,\ A(2\delta n, \delta) \text{ occurs}\}$ (which are δ-separated by definition), and find that $|M(\delta)| \geq O(\delta^-2)$. This allows us to conclude that the Minkowski dimension is two. We summarize these thoughts in the following theorem.

Theorem 4.12. *The tree $\mathcal{T}$ is comapct and has Minkowski dimension two.*

4.4 Convergence with respect to the GH topology

As $\mathcal{T}$ is defined via a limit as $t \to \infty$, and the uniform tree has a paramter $n \to \infty$ we need a technical lemma that takes control of the double limit.

Lemma 4.13. *Let (X, d) be a metric space. Let $(X_n(t) : t \geq 0)_{n=1}^{\infty}$ be a sequence of random functions $[0, \infty) \to X$ and let $(X_n)_{n=1}^{\infty}$ be a sequence of random variables taking values in X. Suppose, for each fixed $t \geq 0$ there exists some random variable $Y(t)$ such that $X_n(t) \xrightarrow{\mathcal{L}} Y(t)$ as $n \to \infty$,and further that there is some random variable Y such that $Y(t) \xrightarrow{\mathcal{L}} Y$ as $t \to \infty$. If*

$$\lim_{t\to\infty} \limsup_{n\to\infty} \mathbf{P}(d(X_n(t), X_n) \geq \epsilon) = 0$$

for all $\epsilon > 0$, then $X_n \xrightarrow{\mathcal{L}} Y$ as $n \to \infty$.

This result is a consequence of the Portmanteau lemma [Bil13]. There are some more technical lemmas, but first let's set up the structures for the actual proof. In particular we need a "continuum" version of a uniform random tree, which is where the discrete algorithm comes in. Let $(U_i)_{i=2}^{n}$ be defined as in the discrete algorithm. We introduce a sequence of random variables $(J_i)_{i=2}^{n}$ with $J_1 = 1$ defined recursively as

$$J_i = (J_{i-1} + 1)\mathbf{1}_{\{U_i < i-1\}} + J_{i-1}\mathbf{1}_{\{U_i \geq i-1\}}.$$

Then, we define a new sequence $(V_i^*)_{i=2}^{n}$ with $V_1 = 0$ in ℓ_1 recursively as

$$V_i^* = V_{\min(U_i, i-1)}^* + z_{J_i}.$$

We define a version of the uniform random tree in ℓ_1 for all $t \geq 0$, denoted by $T_n(t)$ to be the smallest connected set in ℓ_1 containing the set

$$\left\{V_1^*, \ldots, V_{\lfloor t \rfloor \wedge n}^*\right\}.$$

We define $T_n := T_n(n)$. We define similarly to before, $D_n(i, j) = \min_{1\leq m\leq i} \|V_j^* - V_m^*\|$. A result analogous tothe lemma about $D(s, t)$ can be proved about D_n.

Lemma 4.14. *There exists some $K \geq 0$ and $C > 0$ such that for all $k \geq K$,*

$$\mathbf{P}\left(\max_{1\leq m\leq n} D_n(\lfloor\sqrt{n}e^k\rfloor, m) \geq 6ke^{-k}\sqrt{n}\right) \leq Ce^{-k}.$$

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